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Ideal Cost-Free Distributions in Structured Populations for General Payoff Functions

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Abstract The important biological problem of how groups of animals should allocate themselves between different habitats has been modelled extensively. Such habitat selection models have usually involved infinite well-mixed populations. In particular, the model of allocation over a number of food patches when movement is not costly, the ideal free distribution (IFD) model, is well developed. Here we generalize (and solve) a habitat selection game for a finite structured population. We show that habitat selection in such a structured population can have multiple stable solutions (in contrast to the equivalent IFD model where the solution is unique). We also define and study a “predator dilution game” where unlike in the habitat selection game, individuals prefer to aggregate (to avoid being caught by predators due to the dilution effect) and show that this model has a unique solution when movement is unrestricted.

Keywords Structured populations · IFD · Multi-player games · Best response dynamics

1 Introduction

Animals of many types interact in groups in a number of ways [21]. These might be long-standing social groups, as in the case of many species of primate [41], or packs of wild dogs [45, 46]. Alternatively, they might be more transitional in nature, for instance foragers gathering at a food patch, see Krijger and Sevenster [22] for the case of insects and Steele and Hockey [39] for the case of birds. In long-standing groups, there can be complex social relationships and the individuals often form themselves into dominance hierarchies. In more

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transitional groups, relationships will generally be simpler if the members of the group have no prior information about the others. This is the case for the foraging examples that we mentioned above, but also in the case of egg laying as in Prokopy and Roitberg [36]. It is this more transitional type of interaction that is the focus of this paper.

For patch foraging, assuming that individuals are free to move between any of the available food patches, at effectively no cost, animals should follow the ideal free distribution (IFD) over the patches. The simplest single species model was developed by Fretwell and Lucas [14], Fretwell [13] (see also [10]), with a closely related model developed by Parker [35]. More complex cases, including for multi-species [15,23], individuals with different abilities [19,40] and for more complex payoff functions, including the Allee effect [12,25], have also been considered.

It has been shown [9,11,23] that the IFD is an evolutionarily stable strategy [30]. In all of the models above, the population has been infinite, and all foraging patches are reachable by all individuals. Real populations are finite, of course (although large populations can be well modelled as being infinite), but they generally also contain structure. For instance in the case of the wild dogs of Woodroffe et al [45], individuals will only be able to access a subset of the available places (at least without paying a significant cost). Recently, models of populations which are not well mixed, but have some structure, have been developed to take this into account. Evolutionary graph theory models [1,2,5,26–28,42,43] embed standard games such as the prisoner's dilemma within a graph structure [17,34,37]. The interactions within such a structure are restricted to be pairwise, however, and so are unsuited to the type of aggregated situations that we envisage. A more general framework of interactions of populations with restricted movement was introduced in Broom and Rychtář [6]. In that paper, a finite population of individuals distributed themselves over available places subject to their individual restrictions, leading to interactions within the groups on those places, with rewards depending upon the composition of the group, and potentially also on the composition of the groups at other places. More generally the paper introduced a new framework for modelling the evolution of structured populations using multiplayer games (see also [4,8]). In the current paper, we build upon this model by analysing an important class of models, where an individual's movements are strategically chosen and its rewards depend only upon the size of its own group.

2 The Territorial Movement Game

2.1 Game Setting

We consider a population of N individuals I_1, \dots, I_N who can move between and potentially interact on a number of M distinct places P_1, \dots, P_M . The scenario can be illustrated as bipartite graphs as in Fig. 1 where a link connects an individual I_n and a place P_m if and only if the place P_m can be visited by individual I_n [6]. A collection of all places P_m that can be visited by an individual I_n will be called the *territory* of I_n and denoted by $\mathcal{P}_n = \{m; I_n \text{ can go to } P_m\}$. Let \mathcal{I}_m denote the set of individuals that can go to place P_m , i.e. $\mathcal{I}_m = \{n; I_n \text{ can go to } P_m\}$, which we shall call the (potential) *visitors* to P_m . If every individual can go everywhere, we will call the population *unstructured*.

The individuals play the game as follows. Each individual I_n chooses a place P_m in its territory to move to. All individuals make the choice at the same time or, equivalently, without knowledge of what others do. Once the individual has made its choice and visits the place

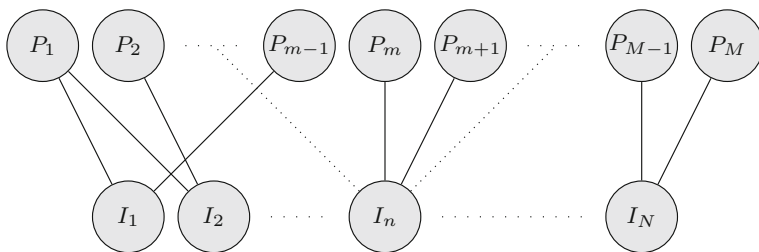


Fig. 1 General situation represented as a bipartite graph. A link connects an individual I_n and a place P_m if and only if the place P_m can be visited by individual I_n

P_m , it receives a reward $R_m(k_m)$ (sometimes later denoted just by R_m) that, for the purpose of this paper, depends solely on the place, P_m , and the (total) number of individuals, k_m , occupying the place. Specifically, the reward does not depend on which individuals share the place and how the other individuals are distributed on the other places; we shall call this the *local aggregation* assumption.

Unpacking this assumption, we are assuming the following. Places can be different of differing quality and of differing accessibility to individuals. Individuals, however, can differ only in the fact that different individuals can (potentially) access different places. This means that once on a particular place, which individuals are there does not affect the payoffs, except through the size of the group. Individuals are only concerned about their immediate payoff, i.e. about what is happening at their current place (and not about what is/ or is not happening at other places).

The strategic element of the game involves the movement of the different individuals; each can employ different moving strategies and their aim is to maximize their payoff.

2.2 Game Theoretical Definitions

In the setting of the Territorial Movement Game, we will use the following definitions which are just the appropriate modification of the usual game theoretical notions (see for an example [7]).

An individual *pure strategy* is the choice of a particular place in a territory.

Given the territorial restrictions, different individuals may have different sets of available strategies to employ.

The *occupancy function* is a function $O : \{1, 2, \dots, N\} \rightarrow \{1, 2, \dots, M\}$ where $P_{O(n)}$ for $O(n) \in \mathcal{P}_n$ is a place to which an individual I_n has moved. For every such occupancy function, we define its *distribution function* as an ordered M -tuple $\langle k_m \rangle_{m=1}^M$ with $k_m = |\{n; O(n) = m\}|$ being the number of individuals at place P_m .

We say that two occupancy functions O_1 and O_2 are *one move apart* (by an individual I_{n_0}) if $O_1(n) = O_2(n)$ for all $n \neq n_0$ while $O_1(n_0) \neq O_2(n_0)$, i.e. if the difference between the occupancies is exactly in one individual changing its position.

The occupancy function O is termed *stable* if any individual would receive a lower payoff when it unilaterally changes its position (given the fixed positions of other individuals); i.e. if for every n

$$R_{O(n)}(k_{O(n)}) > R_m(k_m + 1), \text{ for all } m \in \mathcal{P}_n \setminus O(n). \quad (1)$$

The set \mathcal{O} of occupancy functions is termed a *stable set* if it is a maximal set with the following properties

- (i) if $O, O' \in \mathcal{O}$, then there is a sequence $O = O_1, O_2, \dots, O_k = O'$ of elements in \mathcal{O} such that each pair O_i and O_{i+1} is one move apart.
- (ii) for all $O, O' \in \mathcal{O}$ that are one move apart by an individual I_{n_0} , $R_{O(n_0)}(k_{O(n_0)}) = R_{O'(n_0)}(k'_{O'(n_0)})$,
- (iii) for all $O \in \mathcal{O}$ and $O' \notin \mathcal{O}$ that are one move apart, there must be at least one individual I_{n_0} such that O and O' are one move apart by I_{n_0} and $R_{O(n_0)}(k_{O(n_0)}) > R_{O'(n_0)}(k'_{O'(n_0)})$.

In general it does not make sense to talk about a stable distribution function, as opposed to a stable occupancy function. Two occupancy functions can have the same distribution function, but one may be stable and the other not, because the possible moves depend upon the sets of territories of the individuals at each of the places, which will be different for each occupancy function. The exception is for the well-mixed population, where the territory of each individual is the full set of places, and consequently if an occupancy function is stable, so is any other occupancy function with the same distribution. Thus for the well-mixed population we say that a distribution function is *stable* if any occupancy function with that distribution function (and hence all of them) is (are) stable.

Often when considering evolutionary game theory, we assume that we have so-called *generic games* (see for example [3, 7]). Essentially, assuming that payoffs are derived from the natural world with its underlying variation, we can treat each payoff as if it is the realisation of a random variable, and so we ignore special sets of parameters that would occur with probability 0. This is useful, as the most problematic mathematical cases often occur on such parameter sets. In particular, equalities of distinct payoffs can cause problems. Here we shall call a game *generic* if

$$R_m(k_m) = R_{m'}(k'_m), \text{ if and only if } k_m = k'_m \text{ and } m = m' \quad (2)$$

i.e. if different patches yield different payoffs (irrespective of the group size) and the same patch yields different payoffs for different group sizes. In subsequent sections, we give results that hold for both generic and non-generic games where we can, but sometimes it has proved necessary to restrict ourselves to consideration of generic games only. We clearly indicate at the appropriate point whether generic games are assumed.

A stable occupancy function is always a stable set (of one element). If payoffs are constant and the population is unstructured, the set of all occupancies is a stable set. If the payoffs are generic, a stable set has only one element, and this element is a stable occupancy function.

In Sect. 2.3, we shall see that there is always at least one stable occupancy function for our game, providing that the local aggregation assumption holds. This is not generally true if it does not hold, as we see in the following example.

Example 1 A stable occupancy function may not exist without the local aggregation assumption. Consider two individuals, a predator and its prey, both of whom can go to one of two available places. The first individual, the predator, wants to be with the prey, but the prey wants to avoid the predator. The predator receives payoff 1 (0) if it is (not) at the same place as the prey, the prey receives 0 (1) if it is (not) at the same place as the predator. Note that the payoffs depend on who else is on the place, not only on the number of occupants. Clearly, no pair of pure strategies form a stable occupancy function.

2.3 Existence of a Stable set for the Territorial Movement Game

Here we show that there is always at least one stable set for the Territorial Movement Game. We will consider two Markov processes on the set of all occupancy functions. For the *strong* Markov process, we will consider a transition from occupancy O (with a distribution function

$\langle k_m \rangle_{m=1}^M$) to an occupancy O' (with a distribution function $\langle k'_m \rangle_{m=1}^M$) to be allowable if they are one move apart by n_0 and the individual I_{n_0} would improve its payoff. Thus, there is a potential transition if and only if there are $m_1, m_2 \in \{1, \dots, M\}$ such that

- (i) $\mathcal{I}_{m_1} \cap \mathcal{I}_{m_2} \neq \emptyset$, i.e. there is an individual that could go to both P_{m_1} and P_{m_2} ,
- (ii) $k'_{m_1} = k_{m_1} - 1$, i.e. the individual is (potentially) leaving P_{m_1} ,
- (iii) $k'_{m_2} = k_{m_2} + 1$, i.e. the individual is (potentially) coming to P_{m_2} ,
- (iv) $R_{m_1}(k_{m_1}) < R_{m_2}(k'_{m_2})$, i.e. the payoff the individual gets at P_{m_2} is higher than the one it got at P_{m_1} and
- (v) $k'_m = k_m$ for all $m \neq m_1, m_2$, i.e. nobody else has changed their position.

Note that for the strong process, there cannot be transition from the state O to O' and from O' to O for any pair of states. Furthermore in the generic case, there is no potential transition under the strong process if and only if O is stable.

For the *weak* Markov process, we will consider a transition to be allowable if the conditions (i)–(v) are satisfied, but the inequality in (iv) is not strict. Note that in either process we do not specify the probability of a transition occurring, except that any allowable transition has probability greater than zero.

We shall show the existence of a stable set in two stages.

(a) Let us first define \mathcal{A}_s to be the set of all absorbing states of the strong Markov process. We claim that $\mathcal{A}_s \neq \emptyset$.

Suppose that $\mathcal{A}_s = \emptyset$. Then we could easily find a cycle $O_1, O_2, O_3, \dots, O_T = O_1$ such that transitions can occur from O_i to O_{i+1} in the strong Markov process. This yields the existence of the sequences $\langle m_{1,t} \rangle_{t=1}^{T-1}$ and $\langle m_{2,t+1} \rangle_{t=1}^{T-1}$ such that

$$R_{m_{1,t}}(k_{m_{1,t},t}) < R_{m_{2,t+1}}(k_{m_{2,t+1},t+1}) \quad (3)$$

corresponding to the fact that at time t there were $k_{m_{1,t},t}$ individuals in place $P_{m_{1,t}}$ and one of them was able to improve its position by moving to place $P_{m_{2,t+1}}$. For every transition at step t , we will have one inequality (3). Let us now look at the sequence $\langle k_{m,t} \rangle_{t=1}^T$ for a specific m . The sequence is generally constant with occasional jumps by ± 1 . Any time an individual leaves a place P_m , we have $k_{m,t} > k_{m,t+1} = k_{m,t} - 1$, and the term $R_m(k_{m,t})$ appears on the left-hand side of an inequality (3). Similarly, any time an individual moved to place P_m , $k_{m,t} < k_{m,t+1} = k_{m,t} + 1$, the term $R_m(k_{m,t+1})$ appears on the right-hand side of an inequality (3). Since the sequence $\langle k_{m,t} \rangle_{t=1}^T$ begins and ends at $k_{m,1}$, the term $R_m(k_{m,t})$ has to appear on the left-hand side of (3) as many times as on the right-hand side of (3). So adding those inequalities together (for all t) yields a contradiction.

So, we have shown that there is at least one occupancy $O \in \mathcal{A}_s$. If $O, O' \in \mathcal{A}_s$ are one move apart by an individual I_{n_0} , there is no transition from O to O' or from O' to O in the strong Markov process and thus condition (ii) for a stable set is satisfied for all pairs of elements in \mathcal{A}_s . Consequently, when $O, O' \in \mathcal{A}_s$, there is a transition (both ways) under the weak Markov process between O and O' .

(b) Now, let $\mathcal{B}_w \subset \mathcal{A}_s$ be the set of occupancies that cannot, after any finite number of transitions in the weak Markov process, reach the set \mathcal{A}_s^c , the complement of \mathcal{A}_s . As above, there must be an occupancy function $O \in \mathcal{B}_w$; otherwise, we would again get a cycle (and although we could now potentially have equalities in (3), we would have at least one strict inequality).

Since $\mathcal{B}_w \subset \mathcal{A}_s$, \mathcal{B}_w satisfies condition (ii) for a stable set. Now assume that it does not satisfy condition (iii), i.e. that there is $O \in \mathcal{B}_w$ and $O' \notin \mathcal{B}_w$, i.e. O and O' are one move by I_{n_0} apart so that $R_{O(n_0)}(k_{O(n_0)}) \leq R_{O'(n_0)}(k'_{O'(n_0)})$. But this means that there is a transition

from O to O' in a weak Markov process, and thus, there is a sequence of transitions from O through O' to an element of \mathcal{A}_s^c , a contradiction with the fact that $O \in \mathcal{B}_w$.

A maximal “connected” component of \mathcal{B}_w (i.e. the set of occupancies which one can get from any element of \mathcal{B}_w by a sequence of changes) is then a stable set.

Remark 1 Here we illustrated with an example the concept of the sets \mathcal{A}_s and \mathcal{B}_w from the proof above. Consider 80 individuals with unrestricted movements that can go to two patches with payoffs $R_1(80) = R_2(80) = 1$ and $R_1(k) = R_2(79 - k) = 0$ for $0 \leq k \leq 79$. Then \mathcal{A}_s corresponds to the set of all states but $\langle 79, 1 \rangle$ or $\langle 1, 79 \rangle$ and \mathcal{B}_w consists of $\langle 80, 0 \rangle$ and $\langle 0, 80 \rangle$.

Whilst the above results hold for generic and non-generic games, the results from here until the end of Sect. 2.3 apply to generic games only.

Remark 2 For the generic territorial movement game, consider the algorithm “start with any occupancy function, and make transitions as in the strong Markov process until no further moves can be made”. As previously noted, for the generic case this yields a stable occupancy function. However, there may be of the order of M^N occupancy functions and thus the algorithm can take a long time to finish. We note that which O' is chosen next will in general depend upon the transition probabilities of the strong Markov process, which would need to be specified.

Unfortunately, it is not the case that the maximum improvement decreases every step and it cannot thus be used as an indicator of the fact that the process is already converging to the stable state. For example, consider a population of two individuals. Individual I_1 , currently at P_1 can go to P_1 and P_2 , individual I_2 , currently at P_2 can go to P_2 and P_3 . Let the payoffs be $R_1(1) = 0.5$, $R_2(1) = 1$, $R_2(2) = 0.51$, $R_3(1) = 0.99$. The only improvement that can be made is that individual I_1 moves from place P_1 to place P_2 , increasing its payoff by 0.01. That move, however, sharply decreases the payoff of individual I_2 from $R_2(1) = 1$ to $R_2(2) = 0.51$. The next move would be by individual I_2 from P_2 to P_3 , increasing its payoff by 0.48 to $R_3(1) = 0.99$.

Remark 3 An alternative way to show the existence of a stable occupancy function in the special case of the predator dilution game is shown in Sect. 4.3.

We have shown here that there is at least one stable occupancy function, but in general there may be many, as we see in the following example.

Example 2 Consider a game where individuals prefer to be in groups of even numbers, and the payoff to an individual is 1 (0) if it is (not) in a group of even-numbered size. Let there be $2n$ individuals and m places such that all individuals can go to all places. The number of distributions of individuals over the places is $\binom{2n+m-1}{m-1}$ (see [7] Exercise 9.1). The number of stable distributions can be obtained by sending individuals to places in pairs, yielding $\binom{n+m-1}{m-1}$ stable states. If $n \gg m$, then the number of distributions is roughly $(2n)^{m-1}/((m-1)!)$ and the number of stable distributions is $(n)^{m-1}/((m-1)!)$ so the proportion of stable distributions is 2^{1-m} .

When $m = 2$, the proportion of stable distributions is exactly $(n+1)/(2n+1)$. Thus we can have two-thirds of distributions as stable if we have two individuals on two patches. We conjecture that two-thirds is the largest possible proportion of stable distributions. Note that in this example the proportion of stable *occupancy functions* for $n = 1, m = 2$ is actually $1/2$, since the two stable distributions are each associated with a single occupancy function, but the single unstable distribution is associated with two.

3 A Habitat Selection Game

3.1 Game Setting

Consider M places P_1, P_2, \dots, P_M and N individuals I_1, \dots, I_N . If there are k_m individuals at place P_m , the payoff to each of these individuals will be

$$R_m(k_m) = B_m - f_m(k_m) \quad (4)$$

where $B_m > 0$ is the basic suitability of place P_m and f_m is a nonnegative, increasing and continuous function, with $f_m(0) = 0$. This corresponds to selecting a habitat to forage on, and as the number of individuals on a given patch increases, the actual yield of the patch decreases. However, the yield does not depend on the particular occupants of the patch. In the unstructured population case where every individual can go everywhere, the game has been described for example in Cressman et al. [10], see also Křivan et al. [23].

3.2 A Unique Distribution in an Unstructured Population

The following reasoning, up until the end of Sect. 3, applies to both generic and non-generic games; note, however, that as shown in Remark 4 below, we do not necessarily have uniqueness in non-generic cases. A known solution to this game for effectively infinite populations is the so-called ideal free distribution [14] (ideal = rational individuals, free = no cost or restriction on movement). Typically, conditional on being at a particular place, the individuals want to be with as few other individuals as possible.

(a) Firstly, we show how to find a solution to the habitat selection game in finite unstructured population. Rather than assuming that individuals all make their choice at the same time, let the individuals choose sequentially. The first individual will choose a place P_{m_1} that yields the highest payoff; i.e.

$$R_{m_1}(1) \geq \max_{m' \neq m_1} \{R_{m'}(1)\}. \quad (5)$$

Clearly, individual I_1 cannot do any better; and if no other individual arrives, we have the stable occupancy function if there is a strict inequality in (5), i.e. a unique place with best reward. If there are k places which satisfy (5), then there will be a stable set including k distinct occupancy functions. In this case, assuming that $N \geq k$, the first k individuals must choose the k solutions to (5), since after each choice, the value of the chosen place will decrease, and no longer yield the highest payoff to subsequent individuals. Now, assume that individuals I_1, I_2, \dots, I_i made their choices in a sequence, so that there are k_m individuals at place P_m and that no individual could do better by unilaterally moving to a different place. The individual I_{i+1} will pick $P_{m_{i+1}}$ such that

$$R_{m_{i+1}}(k_{m_{i+1}} + 1) \geq \max_{m' \neq m_{i+1}} \{R_{m'}(k_{m'} + 1)\}. \quad (6)$$

Clearly, in addition to individual I_{i+1} being unable to increase its reward, none of the other individuals I_j for $j < i + 1$ can increase their reward by unilaterally moving to a different place either. Indeed, if I_{i+1} went to the same place as I_j , then I_j must be at the best possible position it could now be (otherwise, I_{i+1} would have chosen a different one). Also, if I_j is at a different place to I_{i+1} , then it cannot improve its payoff by going anywhere else because the move to any other place had not been optimal for I_j even before I_{i+1} arrived (and now the payoff at $P_{m_{i+1}}$ has become even smaller due to the presence of I_{i+1} , whilst the payoffs

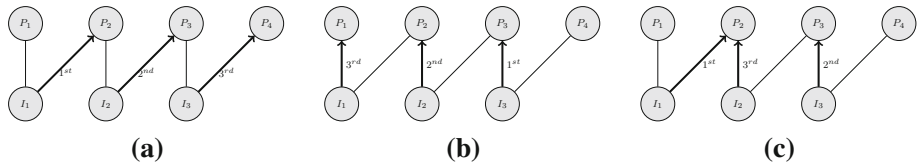


Fig. 2 Method showing the existence for the habitat selection game in an unstructured population cannot be adopted for structured populations. In all instances, we assume that the payoffs are such that $R_2(1) > R_3(1) > R_4(1) > R_1(1) > R_2(2) > R_3(2)$. **a** and **b** show different stable occupancy functions for different orders at which individuals choose their places. **c** Shows that there is an order under which no stable occupancy is reached (one can still follow the procedures from Sect. 2.3 to reach a stable occupancy function in several steps)

at other places have not changed). In this way, we can continue to place the individuals until all have made a choice.

(b) Now we show that while there will usually be many stable occupancy functions (depending on the sequence in which the individuals will make their choice), in unstructured populations and generic games, these will all yield the same distribution $\langle k_m \rangle_{m=1}^M$.

To see the last statement, assume to the contrary that two stable occupancy functions O and O' yield different distributions $\mathbf{k} = \langle k_m \rangle_{m=1}^M$ and $\mathbf{k}' = \langle k'_m \rangle_{m=1}^M$. Because $\mathbf{k} \neq \mathbf{k}'$ and $\sum_m k_m = \sum_m k'_m$, there must be m_1 and m_2 such that $k_{m_1} > k'_{m_1}$ and $k_{m_2} < k'_{m_2}$. Since the payoffs R_{m_1} and R_{m_2} are monotone, we have

$$R_{m_2}(k_{m_2} + 1) \geq R_{m_2}(k'_{m_2}), \quad (7)$$

$$R_{m_1}(k'_{m_1} + 1) \geq R_{m_1}(k_{m_1}). \quad (8)$$

Since O and O' are stable occupancy functions, we have

$$R_{m_1}(k_{m_1}) > R_{m_2}(k_{m_2} + 1), \quad (9)$$

$$R_{m_2}(k'_{m_2}) > R_{m_1}(k'_{m_1} + 1). \quad (10)$$

Thus, by (7), (8), (9), and (10), we get $R_{m_1}(k_{m_1}) > R_{m_1}(k_{m_1})$ which is a contradiction. Consequently, the stable occupancy functions O and O' must yield the same distributions \mathbf{k} and \mathbf{k}' .

Remark 4 In a non-generic game, different occupancies (even in the same stable set) may yield different distributions. For example, consider an unstructured population of three individuals and two places with $R_1(1) > R_1(2) > R_1(3) = R_2(1) > R_2(2) > R_2(3)$. Then occupancies yielding a distribution $\langle 3, 0 \rangle$ and a distribution $\langle 2, 1 \rangle$ are both in the same stable set.

Remark 5 It seems that the same method as presented above would work for the case where individuals could visit only some of the places; the only modification needed being to take the maximum in (5) and (6) only over the allowable places. This is not the case, however, as shown in Fig. 2. Consequently, we have seen that even the well-understood habitat selection game can get more complex if we put restrictions on individual movements.

The fact that the habitat selection game (and Territorial Movement Game in general) has a stable occupancy function no matter what the exact structure was shown in Sect. 2.3. There seems to be no clear argument on how to obtain the solution for a general population in this game, however. This is in contrast to the game that we now consider in Sect. 4.

4 A Predator Dilution Game

4.1 Game Setting

Consider M places P_1, P_2, \dots, P_M and N individuals I_1, \dots, I_N each of which can move to some (but potentially not all) of the places where the payoffs $R_m(k_m)$ are increasing functions of k_m . This corresponds to the following scenario. Each day each individual selects a place. At some point during the day, a predator arrives in the general area, picks a place at random (i.e. with probability $1/M$) and attempts to eat one individual that is on that place. Each individual at the chosen place P_m is equally likely to be eaten, with probability $\lambda_m/(1 + \lambda_m k_m)$, using the Holling type II predation function Křivan [18], where λ_m corresponds to the rate at which the predator finds any given prey (assuming unit available searching time). If there is no individual present, the predator stays hungry. Note that we assume that, unlike in [24], there is no refuge, i.e. $\lambda_m > 0$ for all m .

The individual payoff for staying at place P_m is thus

$$R_m(k_m) = B_m - \frac{1}{M} \frac{\lambda_m}{1 + \lambda_m k_m}, \quad (11)$$

where B_m is a baseline payoff. Many other payoff functions are possible, as long as R_m are all increasing functions of k_m .

In contrast to the habitat selection game, the individuals are now better off if they aggregate (although this would not be the case for a Holling type III function, see for example Garay and Móri [16]). In particular, in the special case where all places are of identical quality and individuals can go to all places, there can be M stable occupancy functions (each with a unique distribution), corresponding to all individuals being at a single place (in contrast to the habitat selection game when different stable occupancy functions yield the same distribution).

Note that if we assumed that the predator keeps searching until it finds an occupied patch, the payoff to an individual could depend on the positions of others and thus violate the local aggregation assumption. For example, consider the case with large N and M with $N \leq M$, and a focal individual being alone. The payoff to the focal individual would be different if the others are all together (it will be killed with probability $1/2$) or if they are all alone (it would be killed with probability $1/N$). For our game defined above, this probability is $1/M$ irrespective of the distribution of the others (as long as they are at different places to the focal individual).

4.2 All Solutions to the Predator Dilution Game

Here we present a method of finding all possible stable occupancy functions of the predator dilution game. The methodology in this section applies to generic games only. The fundamental feature of the payoffs is that they increase with group size. Consequently, the individuals will want to aggregate as much as possible. In other words, they will try to fill up an already occupied place to its maximal capacity.

(a) Let us pick a permutation π on indices $1, 2, \dots, M$ and define the occupancy function O by filling the places to the maximum in the order prescribed by the permutation as follows. We partition $\{1, 2, \dots, N\}$ into a disjoint union of $\{\mathcal{U}_m^I; m = 1, \dots, M\}$ where

$$\mathcal{U}_1^I = \mathcal{I}_{\pi(1)}, \quad (12)$$

$$\mathcal{U}_m^I = \mathcal{I}_{\pi(m)} \setminus \bigcup_{m'=1}^{m-1} \mathcal{I}_{\pi(m')}; 1 < m \leq M. \quad (13)$$

In the above, \mathcal{U}_1^I corresponds to all individuals that can go to $P_{\pi(1)}$ and \mathcal{U}_m^I corresponds to all individuals that can go to $P_{\pi(m)}$ but not to any $P_{\pi(m')}$ for $m' < m$. Define

$$O(n) = \pi(m); \text{ whenever } n \in \mathcal{U}_m^I \quad (14)$$

This yields a distribution $k_{\pi(m)} = |\mathcal{U}_m^I|$. To check for the stability of O , first note that an individual in $P_{\pi(j)}$ cannot physically move to $P_{\pi(i)}$ for $i < j$ (since if $O(n) = \pi(j)$, then $n \in \mathcal{U}_j^I$ and thus $n \notin \mathcal{I}_{\pi(i)}$). Consequently, an occupancy O will be stable if and only if

$$R_{\pi(i)}(k_{\pi(i)}) > R_{\pi(j)}(k_{\pi(j)}) \text{ whenever } i < j \text{ and } \mathcal{I}_{\pi(i)} \cup \mathcal{I}_{\pi(j)} \neq \emptyset. \quad (15)$$

(b) We claim that by going through all the permutations π , identifying the partitions $\{\mathcal{U}_m^I\}$, defining the occupancy functions by (14) and discarding any non-stable ones (those that do not satisfy (15)), we will get all the stable occupancy functions.

First, it is clear that any occupancy function we get by the above procedure is stable. Conversely, if O is a stable occupancy function with the distribution $\langle k_m \rangle_{m=1}^M$, consider a permutation π so that the sequence $\langle R_{\pi(m)}(k_{\pi(m)}) \rangle_{m=1}^M$ is decreasing. The place $P_{\pi(1)}$ must be fully occupied; otherwise, any individual in $\mathcal{I}_{\pi(1)}$ but not currently at $P_{\pi(1)}$ would do better by moving to $P_{\pi(1)}$ since for all m

$$R_{\pi(1)}(k_{\pi(1)} + 1) > R_{\pi(1)}(k_{\pi(1)}) \geq R_m(k_m). \quad (16)$$

Consequently, all individuals with $P_{\pi(1)}$ in their territory must go there. Similarly, the remaining individuals with $P_{\pi(2)}$ in their territory have to go to $P_{\pi(2)}$ and so on. So, the above procedure with the permutation π indeed recovers the stable occupancy function D , and any stable distribution which satisfies (15) must have an occupancy function defined by (12) and (13).

4.3 Existence of Stable Occupancy Functions for the Predator Dilution Game

Following Sect. 2, we know that there is at least one stable set (a stable state for the generic game) for the general territorial movement game, and so in particular for the predator dilution game. In Sect. 4.2, we have seen a way of finding all solutions to the predator dilution game, but there may be few or many, and each case needs to be searched one by one. Here we show a method to efficiently find at least one solution to the predator dilution game (particularly useful for large populations/ structures).

The existence of stable occupancy functions for the Territorial Movement Game has been already shown in Sect. 2.3. However, due to the monotonicity of rewards in the predator dilution game, we can provide an explicit recursive method to identify a stable occupancy function.

We will define a permutation π on $\{1, \dots, M\}$ as follows. At the beginning, consider that all individuals are unplaced and all places unoccupied, and so define the sets of unoccupied places and unplaced individuals at step 0 as

$$\mathcal{V}_0^I = \{1, 2, \dots, N\}, \quad (17)$$

$$\mathcal{V}_0^P = \{1, 2, \dots, M\}. \quad (18)$$

Then in each recursive step, we find a place that gives its occupants maximal reward (from all unoccupied places) when occupied by a maximum number of (yet unplaced) individuals.

This means that once \mathcal{V}_{t-1}^I and \mathcal{V}_{t-1}^P are defined, for some $t \geq 1$, we pick $\pi(t) \in \mathcal{V}_{t-1}^P$ such that

$$R_{\pi(t)}(|\mathcal{I}_{\pi(t)} \cap \mathcal{V}_{t-1}^I|) = \max_{m \in \mathcal{V}_{t-1}^P} R_m(|\mathcal{I}_m \cap \mathcal{V}_{t-1}^I|). \quad (19)$$

In principle, in (19) we may have to evaluate $R_m(0)$ for some m and for simplicity we define it as $R_m(0) = -\infty$. When $\pi(t)$ is chosen, we define

$$\mathcal{V}_t^I = \mathcal{V}_{t-1}^I \setminus \mathcal{I}_{\pi(t)} \text{ and} \quad (20)$$

$$\mathcal{V}_t^P = \mathcal{V}_{t-1}^P \setminus \{\pi(t)\}, \quad (21)$$

corresponding to the fact that now place $P_{\pi(t)}$ is fully occupied and individuals that could go to $P_{\pi(t)}$ are now placed.

We repeat the process until $\mathcal{V}_t^I = \emptyset$. Since we are filling the places to their maximum, $\mathcal{V}_t^P = \emptyset$ implies $\mathcal{V}_t^I = \emptyset$. Also, each step of the above procedure will make exactly one more place occupied and we are also assigning at least one individual to such a place. Consequently, the procedure is well defined and will stop in no more than $\min\{N, M\}$ steps. If the above procedure has not defined π on the whole $\{1, \dots, M\}$, we will just define π to be the (unique) increasing numerical sequence on the remaining elements.

Once the permutation π is defined, we use it to construct the occupancy function O as done earlier in Sect. 4.2 by (14) together with (12) and (13). The occupancy function O is clearly stable. Indeed, the individuals that go to $P_{\pi(1)}$ cannot do any better as this place is fully occupied and it is the place that yields a payoff higher than at any other fully (and thus also arbitrarily) occupied place. The individuals at place $P_{\pi(2)}$ also cannot do any better, because they cannot move to $P_{\pi(1)}$ but are at the place that is fully occupied and yields the highest reward of all other places occupied fully (and thus also arbitrarily) by those that cannot go to $P_{\pi(1)}$; and so on.

5 Discussion

In this paper, we have considered the important problem of how groups of animals should allocate themselves between different habitats, introducing a method to model this using a finite population, where the movement of the individuals faces some restriction. We have considered the general case, showing that there is at least one stable population distribution, but there can be many. We have also considered a development of the classical ideal free distribution of optimal foraging to our structure and developed a model of predator dilution; in each case, we have shown how to find the stable occupancy functions (and that the corresponding distribution is unique in the case of unrestricted movement for the generic case).

We have made no explicit mention of evolution for our model. For evolutionary processes involving finite populations, the specific dynamics involved is generally an important factor, and the most important concept is the fixation probability. A generalisation of the ESS to a finite population, involving a combination of the relative fitness of a single mutant individual when compared to the resident fitness and the fixation probability, was considered in Nowak [33]. Alternatively, and perhaps more appropriately for the current model, we could consider an infinite population of individuals, from which random groups of N individuals are selected to play the above game. In this scenario, strategies that we have termed stable occupancy functions would become ESSs (and similarly stable sets would become evolutionarily stable sets). We note, however, that for any complex structure of places and individuals it is unlikely

that it would be replicated from group to group within the overall population, and so some more complicated scenario with many possible structures might be needed.

This model reflects real cases where foraging is restricted. An example would be pollinators (as the individuals) and plants (as the places), where each pollinator specialises in a subset of the available plants and some plants specialize for certain pollinators [20,44]. Here the “territories” of individuals are not physical areas in space, but their own set of host plants. It can also be used to model more general situations, which are not principally about foraging. A range of activities that apply to group-living animals can be regarded in this way, for example taking shelter, or migration (where a “place” would be a migration route) or alternatively mating, where the females may be considered as individuals and the males as “places”. Currently our model is very general, but it is flexible, and there is the potential to develop it to model more species-specific situations.

There are a number of ways that the modelling in this paper can be developed. One can relax the local aggregation assumption as done for example in Maciejewski and Puleo [29] where authors consider a structured population model where different individuals are best suited to different regions of their environment. Also, in this paper we have assumed strategies are pure, so that each individual chooses a place so that given the individuals choices the occupancy function is determined. An interesting question is, can there be mixed strategies, where individuals choose a place according to a probability distribution? We note that a population with only two individuals moving to M patches is a form of bi-matrix game (with special payoffs). ESSs in such games are always pure, following the work of Selten [38]. When can this result be generalised to our multi-player case, meaning that no mixed strategy could be stable? It would also be of interest to consider general, and specific, models on structures of movement with special properties. Models in evolutionary graph theory have focused on such ideas to significant benefit, and there is potential in our model too.

Further, we note that the habitat selection game and the predator dilution game are effectively two extremes of a more general case, in the sense that the ideal group size for the habitat selection game was one (i.e. with no groupmates) and for the predator dilution game individuals prefer to be with as many others as possible. In reality, the ideal group size may typically be somewhere between these two. For example, if individuals are hunters, then a group has to be large enough to catch a large prey animal but not so large that the share of every individual is overly reduced. Similarly, there may be an optimal group size to hide from predators. If the group is too big, the predator dilution effect is strong, but the group may be very visible to potential predators. If the group is too small, the dilution effect is small if the group is found by a predator. Consequently, the optimal group size is neither small nor large.

Thus it would be useful to consider such a generalisation where the best group size is of intermediate value. This would be appropriate for any population which is subject to the Allee effect. This type of situation was modelled in the usual infinite population habitat selection game in Křivan [25] (see also [31,32]) which demonstrates the strong influence of the Allee effect upon the population dynamics and also upon the spatial distribution of the population. A unique IFD exists for low and high population densities, but there can be up to three IFDs for intermediate densities. Thus even in the case of well-mixed infinite populations, the number of possibilities is large. Thus we do not expect to be able to generate similar complete general results for this case as those we found for the previous two models. Nevertheless we know from Sect. 2.3 that there will always be at least one stable set, and we envisage that some general results can be found, together with approaches to tackle specific models, including finding the stable sets, using the methods we have generated in this paper. It would also be of interest to attempt to combine the analytical methods developed in our paper with those of Křivan [25] for the infinite population unrestricted case.

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