Chapter 1
Biology and Evolutionary Games

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Abstract This chapter surveys some evolutionary games used in biological sciences. These include the Hawk-Dove game, the Prisoner’s Dilemma, Rock–Paper–Scissors, the war of attrition, the Habitat Selection game, predator-prey games, and signalling games.

Key words: Battle of the Sexes, Foraging games, Habitat Selection game, Hawk-Dove game, Prisoner’s Dilemma, Rock–Paper–Scissors, Signalling games, War of Attrition

1.1 Introduction

Evolutionarily game theory (EGT) as conceived by Maynard Smith and Price (1973) was motivated by evolution. Several authors (e. g., Wynne-Edwards,
at that time argued that animal behavior patterns were “for the good of the species” and that natural selection acts at the group level. This point of view was at odds with the Darwinian viewpoint where natural selection operates on the individual level. In particular, adaptive mechanisms that maximize a group benefit do not necessarily maximize individual benefit. This led Maynard Smith and Price (1973) to develop a mathematical model of animal behavior, called the Hawk–Dove game, that clearly shows the difference between group selection and individual selection. We thus start this chapter with the Hawk–Dove game.

Today, evolutionary game theory is one of the milestones of evolutionary ecology as it put the concept of Darwinian evolution on solid mathematical grounds. Evolutionary game theory has spread quickly in behavioral and evolutionary biology with many influential models that change the way that scientists look at evolution today. As evolutionary game theory is non-cooperative, where each individual maximizes its own fitness, it seemed that it cannot explain cooperative or altruistic behavior that was easy to explain on the grounds of the group selection argument. Perhaps the most influential model in this respect is the Prisoner’s Dilemma (Poundstone, 1992), where the evolutionarily stable strategy leads to a collective payoff that is lower than the maximal payoff the two individuals can achieve if they were cooperating. Several models within evolutionary game theory have been developed that show how mutual cooperation can evolve. We discuss some of these models in Section 1.3. A popular game played by human players across the world, which can also be used to model some biological populations, is the Rock–Paper–Scissors game (RPS; Section 1.4). All of these games are single species matrix games, so that their payoffs are linear, with a finite number of strategies. An example of a game that cannot be described by a matrix and that has a continuum of strategies is the war of attrition in Section 1.5.1 (or alternatively the Sir Philip Sidney game mentioned in Section 1.10). A game with non-linear payoffs which examines an important biological phenomenon is the sex-ratio game in Section 1.5.2.

Although evolutionary game theory started with consideration of a single species, it was soon extended to two interacting species. This extension was not straightforward, because the crucial mechanism of a (single-species) EGT, that is negative frequency dependence that stabilizes phenotype frequencies at an equilibrium, is missing if individuals of one species interact with individuals of another species. These games are asymmetric, because the two contestants are in different roles (such asymmetric games also occur within a single species). Such games that can be described by two matrices are called bimatrix games. Representative examples include the Battle of the Sexes (Section 1.6.1) and the owner-intruder game (Section 1.6.2). Animal spatial distribution that is evolutionarily stable is called the Ideal Free Distribution (Section 1.7). We discuss first the IFD for a single species and then for two species. The resulting model is described by four matrices, so it is no longer a bimatrix game. The IFD, as an outcome of animal dispersal, is
related to the question of under which conditions animal dispersal can evolve (Section 1.8). Section 1.9 focuses on foraging games. We discuss two models that use EGT. The first model, that uses decision trees, is used to derive the diet selection model of optimal foraging. This model asks what the optimal diet of a generalist predator is in an environment that has two (or more) prey species. We show that this problem can be solved using the so called agent normal form of an extensive game. We then consider a game between prey that try to avoid their predators and predators that aim to capture prey. The last game we consider in some detail is a signalling game of mate quality, that was developed to help explain the presence of costly ornaments, such as the peacock’s tail.

We conclude with a brief section discussing a few other areas where evolutionary game theory has been applied. However, a large variety of models that use EGT have been developed in the literature, and it is virtually impossible to survey all of them.

1.2 The Hawk-Dove game: Selection at the individual level vs. selection at the population level

One of the first evolutionary games was introduced to understand evolution of aggressiveness among animals (Maynard Smith and Price, 1973). Although many species have strong weapons (e.g., teeth or horns), it is a puzzling observation that in many cases the antagonistic encounters do not result in a fight. In fact, such encounters often result in a complicated series of behaviors, but without causing serious injuries. For example, in contests between two male red deer, the contestants first approach each other and provided one does not withdraw, the contest escalates to a roaring contest and then to the so called parallel walk. Only if this does not lead to the withdrawal of one deer, does a fight follow. It was observed (Maynard Smith, 1982) that out of 50 encounters, only 14 resulted in a fight. The obvious question is why animals do not always end up in a fight? As it is good for an individual to get the resource (in the case of the deer, the resources are females for mating), Darwinian selection seems to suggest that individuals should fight whenever possible. One possible answer why this is not the case is that such a behavior is for the good of the species, because any species following this aggressive strategy would die out quickly. If so, then we should accept that the unit of selection is not an individual, and abandon H. Spencer’s “survival of the fittest” (Spencer, 1864).

The Hawk-Dove model explains animal contest behavior from the Darwinian point of view. The model considers interactions between two individuals from the same population that meet in a pairwise contest. Each individual uses one of the two strategies called Hawk and Dove. An individual playing Hawk is ready to fight when meeting an opponent while an individual
playing Dove does not escalate the conflict. The game is characterized by two positive parameters where $V$ denotes the value of the contested resource and $C$ is the cost of the fight measured as the damage one individual can cause to his opponent. The payoffs for the row player describe the increase/decrease in the player’s fitness after an encounter with an opponent. The game matrix is

$$
\begin{pmatrix}
\text{Hawk} & \text{Dove} \\
\text{Hawk} & \frac{V-C}{2} & V \\
\text{Dove} & 0 & \frac{V}{2}
\end{pmatrix}
$$

and the model predicts that when the cost of the fight is lower than is the reward obtained by getting the resource, $C < V$, all individuals should play the Hawk strategy that is the strict Nash equilibrium (NE) (thus an evolutionarily stable strategy (ESS)) of the game. When the cost of a fight is larger than the reward obtained by getting the resource, $C > V$, then $p = V/C$ ($0 < p < 1$) is the corresponding monomorphic ESS. In other words, each individual will play Hawk when encountering an opponent with probability $p$ and Dove with probability $1 - p$. Thus, the model predicts that aggressiveness in the population decreases with the cost of fighting. In other words, the species that possess strong weapons (e.g., antlers in deers) should solve conflicts with very little fighting.

Can individuals obtain a higher fitness when using a different strategy? In a monomorphic population where all individuals use a mixed strategy $0 < p < 1$, the individual fitness and the average fitness in the population are the same and equal to

$$
E(p, p) = \frac{V}{2} - \frac{C}{2}p^2.
$$

This fitness is maximized for $p = 0$, i.e., when the level of aggressiveness in the population is zero, all individuals play the Dove strategy, and individual fitness equals $V/2$. Thus, if selection operated on a population or a species level, all individuals should be phenotypically Doves who never fight. However, the strategy $p = 0$ cannot be an equilibrium from an evolutionary point of view, because in a Dove only population, Hawks will always have a higher fitness ($V$) than Doves ($V/2$), and will invade. In other words, the Dove strategy is not resistant to invasion by Hawkish individuals. Thus, securing all individuals to play the strategy $D$, that is beneficial from the population point of view, requires some higher organizational level that promotes cooperation between animals (Dugatkin and Reeve, 1998, see also Section 1.3).

On the contrary, at the evolutionarily stable equilibrium $p^* = V/C$, individual fitness

$$
E(p^*, p^*) = \frac{V}{2} \left(1 - \frac{V}{C}\right)
$$

is always lower than $V/2$. However, the ESS strategy cannot be invaded by any other single mutant strategy.
Darwinism assumes that selection operates at the level of an individual, which is then consistent with non-cooperative game theory. However, this is not the only possibility. Some biologists (e.g., Gilpin, 1975) postulated that selection operates on a larger unit, a group (e.g., a population, a species etc.), maximizing the benefit of this unit. This approach was termed group selection. Alternatively, Dawkins (1976) suggested that selection operates on a gene level. The Hawk-Dove game allows us to show clearly the difference between the group and Darwinian selections.

Group selection vs. individual selection also nicely illustrates the so called tragedy of the commons (Hardin, 1968) (based on an example given by the English economist William Forster Lloyd) that predicts deterioration of the environment, measured by the fitness, in an unconstrained situation where each individual maximizes its profit. For example, when a common resource (e.g., fish) is over-harvested, the whole fishery collapses. To maintain a sustainable yield, regulation is needed that prevents over-exploitation (i.e., which does not allow Hawks that would over-exploit the resource to enter). Effectively, such a regulatory body keeps $p$ at zero (or close to it), to maximize the benefits for all fishermen. Without such a regulatory body, Hawks would invade and necessarily decrease the profit for all. In fact, as the cost $C$ increases (due to scarcity of resources), fitness at the ESS decreases and when $C$ equals $V$, fitness is zero.

### 1.2.1 Replicator dynamics for the Hawk-Dove game

In the previous section we have assumed that all individuals play the same strategy, either pure, or mixed. If the strategy is mixed, each individual randomly chooses one of its elementary strategies on any given encounter according to some given probability. In this monomorphic interpretation of the game the population mean strategy coincides with the individual strategy. Now we will consider a distinct situation where $n$ phenotypes exist in the population. In this polymorphic setting we say that a population is in state $p \in \Delta_n$ (where $\Delta_n = \{ p \in \mathbb{R}^n | p_i \geq 0, p_1 + \cdots + p_n = 1 \}$ is a probability simplex) if $p_i$ is the proportion of the population using strategy $i$. As opposed to the monomorphic case, in this polymorphic interpretation the individual strategies and the mean population strategy are different, because the mean strategy characterizes the population, not a single individual.

The ESS definition does not provide us with a mechanistic description of phenotype frequency dynamics that would converge to an ESS. One of the frequency dynamics often used in evolutionary game theory is the replicator dynamics (Taylor and Jonker, 1978). Replicator dynamics assume that the population growth rate of each phenotype is given by its fitness and they focus on changes in phenotypic frequencies in the population (see Volume I, Chapter Evolutionary games). Let us consider the replicator equation for the
Hawk-Dove game. Let $x$ be the frequency of Hawks in the population. The fitness of a Hawk is

$$E(H, x) = \frac{V - C}{2} x + V(1 - x)$$

and, similarly, the fitness of a Dove is

$$E(D, x) = (1 - x) \frac{V}{2}.$$  

Then the average fitness in the population is

$$E(x, x) = xE(H, x) + (1 - x)E(D, x) = \frac{V - Cx^2}{2}$$

and the replicator equation is

$$\frac{dx}{dt} = x (E(H, x) - E(x, x)) = \frac{1}{2} x(1 - x)(V - Cx).$$

Assuming $C > V$, we remark that the interior distribution equilibrium of this equation, $x = V/C$, corresponds to the mixed ESS for the underlying game. In this example phenotypes correspond to elementary strategies of the game. It may be that phenotypes also correspond to mixed strategies.

1.3 The Prisoner’s Dilemma and the evolution of cooperation

The Prisoner’s Dilemma (see Flood, 1952; Poundstone, 1992) is perhaps the most famous game in all of game theory, with applications from areas including economics, biology and psychology. Two players play a game where they can Defect or Cooperate, yielding the payoff matrix

\[
\begin{pmatrix}
\text{Cooperate} & \text{Defect} \\
\text{Cooperate} & R & S \\
\text{Defect} & T & P \\
\end{pmatrix}
\]

These abbreviations are derived from Reward (reward for cooperating), Temptation (temptation for defecting when the other player cooperates), Sucker (paying the cost of cooperation when the other player defects), and Punishment (paying the cost of defecting). The rewards satisfy the conditions $T > R > P > S$. Thus while Cooperate is Pareto efficient (in the sense that it is impossible to make any of the two players better off without making the other player worse off), Defect row dominates Cooperate and so is the unique ESS, even though mutual cooperation would yield the greater payoff. Real
human (and animal) populations, however, involve a lot of cooperation; how is that enforced?

There are many mechanisms for enabling cooperation, see for example Nowak (2006). These can be divided into six types as follows:

1. Kin selection, that occurs when the donor and recipient of some apparently altruistic act are genetic relatives.
2. Direct reciprocity, requiring repeated encounters between two individuals.
3. Indirect reciprocity, based upon reputation. An altruistic individual gains a good reputation, which means in turn that others are more willing to help that individual.
4. Punishment, as a way to enforce cooperation.
5. Network reciprocity, where there is not random mixing in the population and cooperators are more likely to interact with other cooperators.
6. Multi-level selection, alternatively called group selection, where evolution occurs on more than one level.

We discuss some of these concepts below.

1.3.1 Direct reciprocity

Direct reciprocity requires repeated interaction, and can be modelled by the Iterated Prisoner’s Dilemma (IPD). The IPD involves playing the Prisoner’s Dilemma over a (usually large) number of rounds, and thus being able to condition choices in later rounds on what the other player played before. This game was popularised by Axelrod (1981, 1984) who held two tournaments where individuals could submit computer programs to play the IPD. The winner of both tournaments was the simplest program submitted, called Tit for Tat (TFT), which simply cooperates on the first move and then copies its opponent’s previous move.

TFT here has three important properties: it is nice so it never defects first; it is retaliatory so it meets defection with a defection next move; it is forgiving so even after previous defections, it meets cooperation with cooperation next move. TFT effectively has a memory of one place, and it was shown in Axelrod and Hamilton (1981) that TFT can resist invasion by any strategy that is not nice if it can resist both Always Defect (ALLD) and Alternative (ALT), which defects (cooperates) on odd (even) moves. However, this does not mean that TFT is an ESS, because nice strategies can invade by drift as they receive identical payoffs to TFT in a TFT population (Bendorf and Swistak, 1995). We note that TFT is not the only strategy that can promote cooperation in the IPD; others include tit for two tats (TF2T which defects only after two successive defections of its opponent), Grim (which defects on all moves after its opponent’s first defection) and win stay/lose shift (which changes its choice if and only if its opponent defected on the previous move).
Games between TFT, ALLD and ALT against TFT have the following sequence of moves:

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We assume that the number of rounds is not fixed, and that there is always the possibility of a later round (otherwise the game can be solved by backwards induction, yielding ALLD as the unique NE strategy). At each stage there is a further round with probability \( w \) (as in the second computer tournament), the payoffs are then

\[
E(TFT, TFT) = R + Rw + Rw^2 + \ldots = \frac{R}{1 - w},
\]  

\[
E(ALLD, TFT) = T + Pw + Pw^2 + \ldots = T + \frac{Pw}{1 - w},
\]  

\[
E(ALT, TFT) = T + Sw + Tw^2 + Sw^3 + \ldots = \frac{T + Sw}{1 - w^2}.
\]

Thus TFT resists invasion if and only if

\[
\frac{R}{1 - w} > \max \left( \frac{T + Pw}{1 - w}, \frac{T + Sw}{1 - w^2} \right),
\]

i.e., if and only if

\[
w > \max \left( \frac{T - R}{T - P}, \frac{T - R}{R - S} \right),
\]

i.e., when the probability of another contest is sufficiently large (Axelrod, 1981, 1984). We thus see that for cooperation to evolve here, the extra condition \( 2R > S + T \) is required, since otherwise the right hand side of inequality (1.5) would be at least 1.

Whilst TFT proved successful at promoting cooperation above, what if errors occur, so that an intention to cooperate becomes a defection (or is perceived as such)? After a single mistake a pair of interacting TFT players will be locked in an alternating cycle of Defect versus Cooperation, and then mutual defection after a second mistake when C was intended. Under such circumstances TF2T can maintain cooperation, whereas TFT cannot. In fact a better strategy (in the sense that it maintains cooperation when playing against itself but resists invasion from defecting strategies) is GTFT (generous tit for tat; see Komorita et al, 1968), which combines pure cooperation with TFT by cooperating after a cooperation, but meeting a defection with a defection with probability
1.3.2 Kin selection and Hamilton’s rule

In most evolutionary game theoretical models, the aim of each individual is to maximise its own fitness, irrespective of the fitness of other individuals. However, if individuals are related, then the fitnesses of others should be taken into account.

Let us consider two interacting individuals, with coefficient of relatedness $r$, which is the probability that they share a copy of a given allele. For example father and son will have $r = 1/2$. One individual acts as a potential donor, the other as a recipient, which receives a benefit $b$ from the donor at the donor’s cost $c$. The donating individual pays the full cost, but also indirectly receives the benefit $b$ multiplied by the above factor $r$. Thus donation is worthwhile provided that

$$rb > c \text{ i.e. } r > \frac{c}{b}$$

which is known as Hamilton’s rule (Hamilton, 1964).

Note that this condition is analogous to the condition for cooperation to resist invasion in the IPD above, where a commonly used special class of the PD matrix has payoffs representing cooperation as making a donation and defecting as not. Then TFT resists invasion when $w > c/b$.

1.3.3 Indirect reciprocity and punishment

The IPD is an example of Direct reciprocity. Suppose now we have a population of individuals who play many contests, but these are not in long sequences against a single “opponent” as above? If faced with a series of single shot games, how can cooperation be achieved?

Such situations are often investigated by the use of public goods games involving experiments with groups of real people, as in the work of Fehr and Gachter (2002). In these experiments individuals play a series of games, each game involving a new group. In each game there were four individuals, each of them receiving an initial endowment of 20 dollars, and each had to choose a level of investment into a common pool. Any money that was invested increased by a factor of 1.6 and was then shared between the four individuals, meaning that the return for each dollar invested was 40 cents to each of the players. In particular the individual making the investment of one dollar only receives 40 cents, and so makes a loss of 60 cents. Thus, like the Prisoner’s Dilemma, it is clear that the best strategy is to make
no investment, but simply to share rewards from the investments of other players. In these experiments investment levels began reasonably high, but slowly declined, as players saw others cheat.

In later experiments, each game was played over two rounds, an investment round and a punishment round, where players were allowed to punish others. In particular every dollar “invested” in punishment levied a fine of three dollars on the target of the punishment. This led to investments which increased from their initial level, as punishment brought cheating individuals into line. It should be noted that in a population of individuals many, but not all of whom, punish, optimal play for individuals in this case should not be to punish, but to be a second-order free rider who invests but does not punish, and therefore saves the punishment fee. Such a population would collapse down to no investment after some number of rounds. Thus it is clear that the people in the experiments were not behaving completely rationally.

Thus we could develop the game to have repeated rounds of punishment. An aggressive punishing strategy would then in round 1, punish all defectors; in round 2, punish all cooperators who did not punish defectors in round 1; in round 3, punish all cooperators who did not punish in round 2 as above, and so on. Thus such players not only punish cheats, but anyone who does not play exactly as they do. Imagine a group of $m$ individuals with $k$ cooperators (who invest and punish), $\ell$ defectors and $m - k - \ell - 1$ investors (who do not punish). This game, with this available set of strategies, requires two rounds of punishment as described above. The rewards to our focal individual in this case will be

$$R = \begin{cases} 
\frac{(m-\ell)cV}{m} - kP & \text{if an investor,} \\
\frac{(m-\ell)cV}{m} - (m - k - 1) & \text{if a cooperator,} \\
\frac{(m-\ell-1)cV}{m} + V - kP & \text{if a defector,} 
\end{cases}$$

where $V$ is the initial level of resources of each individual, $c < m$ is the return on investment (every dollar becomes $1 + c$ dollars), and $P$ is the punishment multiple (every dollar invested in punishment generates a fine of $P$ dollars).

The optimal play for our focal individual is

- Defect if $V \left(1 - \frac{c}{m}\right) > kP - (m - k - 1)$,
- Cooperate otherwise.

Thus defect is always stable and invest and punish is stable if $V(1 - c/m) < (m - 1)P$.

We note that there are still issues on how such punishment can emerge in the first place (Sigmund, 2007).


1.4 The Rock–Paper–Scissors game

The Rock–Paper–Scissors game is a three strategy matrix game, which people commonly play recreationally. In human competition the game dates back at least to 17th century China. There is a lot of potential psychology involved in playing the game, and there are numerous tournaments involving it. The important feature of the game is that Rock beats Scissors, Scissors beats Paper and Paper beats Rock. The payoff matrix is

\[
\begin{pmatrix}
\text{Rock} & \text{Scissors} & \text{Paper} \\
0 & a_3 & -b_2 \\
-b_3 & 0 & a_1 \\
a_2 & -b_1 & 0 \\
\end{pmatrix},
\]

where all \(a\)'s and \(b\)'s are positive. For the conventional game played between people \(a_i = b_i = 1\) for \(i = 1, 2, 3\).

There is a unique internal NE of the above game given by the vector

\[
p = \frac{1}{K} (a_1a_3 + b_1b_2 + a_1b_1, a_1a_2 + b_2b_3 + a_2b_2, a_2a_3 + b_1b_3 + a_3b_3),
\]

where the constant \(K\) is just the sum of the three terms to ensure that \(p\) is a probability vector. In addition, \(p\) is a globally asymptotically stable equilibrium of the replicator dynamics if and only if \(a_1a_2a_3 > b_1b_2b_3\). It is an ESS if and only if \(a_1 - b_1, a_2 - b_2\) and \(a_3 - b_3\) are all positive, and the largest of their square roots is smaller than the sum of the other two square roots (Hofbauer and Sigmund, 1998). Thus if \(p\) is an ESS of the RPS game, then it is globally asymptotically stable under the replicator dynamics. However, since the converse is not true, the RPS game provides an example illustrating that whilst all internal ESSs are global attractors of the replicator dynamics, not all global attractors are ESSs.

We note that the case when \(a_1a_2a_3 = b_1b_2b_3\) (including the conventional game with \(a_i = b_i = 1\)) leads to closed orbits of the replicator dynamics, and a stable (but not asymptotically stable) internal equilibrium. This is an example of a non-generic game, where minor perturbations of the parameter values can lead to large changes in the nature of the game solution.

This game is a good representation for a number of real populations. The most well known of these is among the common side-blotched lizard \(Uta stansburiana\). This lizard has three types of distinctive throat colouration, which correspond to very different types of behaviour. Males with orange throats are very aggressive and have large territories which they defend against intruders. Males with dark blue throats are less aggressive, and hold smaller territories. Males with yellow stripes do not have a territory at all, but bear a strong resemblance to females and use a sneaky mating strategy. It was observed in Sinervo and Lively (1996) that if the Blue strategy is the
most prevalent Orange can invade, if Yellow is prevalent Blue can invade and if Orange is prevalent then Yellow can invade.

An alternative real scenario is that of *Escherichia coli* bacteria, involving three strains of bacteria (Kerr et al, 2002). One strain produces the antibiotic colicin. This strain is immune to it, as is a second strain, but the third is not. When only the first two strains are present, the second strain outcompetes the first, since it forgoes the cost of colicin production. Similarly the third outcompetes the second, as it forgoes costly immunity, which without the first strain is unnecessary. Finally, the first strain outcompetes the third, as the latter has no immunity to the colicin.

### 1.5 Non-matrix games

We have seen that matrix games involve a finite number of strategies with a payoff function that is linear in the strategy of both the focal player and that of the population. This leads to a number of important simplifying results (see Volume I, Chapter Evolutionary games). All of the ESSs of a matrix can be found in a straightforward way using the procedure of Haigh (1975).

Further, adding a constant to all entries in a column of a payoff matrix leaves the collection of ESSs (and the trajectories of the replicator dynamics) of the matrix unchanged. Haigh’s procedure can potentially be shortened, using the important Bishop-Cannings Theorem (Bishop and Cannings, 1976), a consequence of which is that if $p_1$ is an ESS, no strategy $p_2$ whose support is either a superset or a subset of the support of $p_1$ can be an ESS.

However, there are a number of ways that games can involve nonlinear payoff functions. Firstly, playing the field games yield payoffs that are linear in the focal player but not in the population (for examples see Sections 1.7.1, 1.9.1). Another way this can happen is to have individual games of the matrix type, but where opponents are not selected with equal probability from the population, for instance if there is some spatial element. Thirdly, the payoffs can be nonlinear in both components. Here strategies do not refer to a probabilistic mix of pure strategies, but a unique trait, such as the height of a tree as in Kokko (2007) or a volume of sperm, see e.g., Ball and Parker (2007). This happens in particular in the context of adaptive dynamics (see Volume I, Chapter Evolutionary games).

Alternatively a non-matrix game can involve linear payoffs, but this time with a continuum of strategies (we note that the cases with nonlinear payoffs above can also involve such a continuum, especially the third type). A classical example of this is the war of attrition (Maynard Smith, 1974).
1.5.1 The war of attrition

We consider a Hawk-Dove game, where both individuals play Dove, but that instead of the reward being allocated instantly, they become involved in a potentially long displaying contest where a winner will be decided by one player conceding, and there is a cost proportional to the length of the contest. An individual’s strategy is thus the length of time it is prepared to wait. Pure strategies are all values of \( t \) on the non-negative part of the real line, and mixed strategies are corresponding probability distributions. These kinds of contests are for example observed in dung flies (Parker and Thompson, 1980).

Choosing the cost to be simply the length of time spent, the payoff for a game between two pure strategies \( S_t \) (wait until time \( t \)) and \( S_s \) (wait until time \( s \)) for the player that uses strategy \( S_t \) is

\[
E(S_t, S_s) = \begin{cases} 
  V - s & t > s, \\
  V/2 - t & t = s, \\
  -t & t < s.
\end{cases}
\]

and the corresponding payoff from a game involving two mixed strategists playing the probability distributions \( f(t) \) and \( g(s) \) to the \( f(t) \) player is

\[
\int_0^\infty \int_0^\infty f(t)g(s)E(S_t, S_s)dtds.
\]

It is clear that no pure strategy can be an ESS, since \( S_s \) is invaded by \( S_t \) (i.e., \( E(S_t, S_s) > E(S_s, S_s) \)) for any \( t > s \), or any positive \( t < s - V/2 \). There is a unique ESS which is found by first considering (analogous to the Bishop-Cannings Theorem; see Volume I, Chapter Evolutionary games) a probability distribution \( p(s) \) that gives equal payoffs to all pure strategies that could be played by an opponent. This is required, since otherwise some potential invading strategies could do better than others, and since \( p(s) \) is simply a weighted average of such strategies, it would then be invaded by at least one type of opponent. Payoff of a pure strategy \( S_t \) played against a mixed strategy \( S_{p(s)} \) given by a probability distribution \( p(s) \) over the time interval is

\[
E(S_t, S_{p(s)}) = \int_0^t (V - s)p(s)ds + \int_t^\infty (-t)p(s)ds. \tag{1.6}
\]

Differentiating equation (1.6) with respect to \( t \) (assuming that such a derivative exists) gives

\[
(V - t)p(t) - \int_t^\infty p(s)ds + tp(t) = 0. \tag{1.7}
\]
If $P(t)$ is the associated distribution function, so that $p(t) = P'(t)$ for all $t \geq 0$, then equation (1.7) becomes

$$VP'(t) + P(t) - 1 = 0$$

and we obtain

$$p(t) = \frac{1}{V} \exp \left( -\frac{t}{V} \right). \quad (1.8)$$

It should be noted that we have glossed over certain issues in the above, for example consideration of strategies without full support or with atoms of probability. This is discussed in more detail in Broom and Rychtar (2013). The above solution was shown to be an ESS in Bishop and Cannings (1976).

### 1.5.2 The sex-ratio game

Why is it that the sex ratio in most animals is close to a half? This was the first problem to be considered using evolutionary game theory (Hamilton, 1967), and its consideration, including the essential nature of the solution, dates right back to Darwin (1871). To maximize the overall birth rate of the species, in most animals there should be far more females than males, given that females usually make a much more significant investment in bringing up offspring than males. This, as mentioned before, is the wrong perspective, and we need to consider the problem from the viewpoint of the individual.

Suppose that in a given population, an individual female will have a fixed number of offspring, but that she can allocate the proportion of these that are male. This proportion is thus the strategy of our individual. As each female (irrespective of its strategy) has the same number of offspring, this number does not help us in deciding which strategy is the best. The effect of a given strategy can be measured as the number of grandchildren of the focal female. Assume that the number of individuals in a large population in the next generation is $N_1$ and in the following generation is $N_2$. Further assume that all other females in the population play the strategy $m$, and that our focal individual plays strategy $p$.

As $N_1$ is large, the total number of males in the next generation is $mN_1$ and so the total number of females is $(1 - m)N_1$. We shall assume that all females (males) are equally likely to be the mother (father) of any particular member of the following generation of $N_2$ individuals. This means that a female offspring will be the mother of $N_2 / ((1 - m)N_1)$ of the following generation of $N_2$ individuals, and a male offspring will be the father of $N_2 / (mN_1)$ of these individuals. Thus our focal individual will have the following number of grandchildren.
To find the best $p$, we maximise $E(p, m)$. For $m < 0.5$ the best response is $p = 1$, and for $m > 0.5$ we obtain $p = 0$. Thus $m = 0.5$ is the interior NE at which all values of $p$ obtain the same payoff. This NE satisfies the stability condition in that $E(0.5, m') - E(m', m') > 0$ for all $m' \neq 0.5$ (Broom and Rychtar, 2013).

Thus from the individual perspective, it is best to have half your offspring as male. In real populations it is often the case that relatively few males are the parents of many individuals, for instance in social groups often only the dominant male fathers offspring. Sometimes other males are actually excluded from the group; lion prides generally consist of a number of females, but only one or two males, for example. From a group perspective these extra males perform no function, but there is a chance that any male will become the father of many.

### 1.6 Asymmetric games

The games we have considered above all involve populations of identical individuals. What if individuals are not identical? Maynard Smith and Parker (1976) considered two main types of difference between individuals. The first type was correlated asymmetries where there were real differences between them, for instance in strength or need for resources, which would mean their probability of success, cost levels, valuation of rewards, set of available strategies etc. may be different, i.e. the payoffs “correlate” with the type of the player. An example of such games are the predator-prey games of Section 1.9.2, or the Battle of the Sexes below in Section 1.6.1.

The second type, uncorrelated asymmetries, occurred when the individuals were physically identical, but nevertheless occupied different roles, for example one was the owner of a territory and the other was an intruder, which we shall see in Section 1.6.2. For uncorrelated asymmetries, even though individuals do not have different payoff matrices, it is possible to base their strategy upon the role that they occupy. As we shall see, this completely changes the character of the solutions that we obtain.

We note that the allocation of distinct roles can apply to games in general, for example there has been significant work on the asymmetric war of attrition (see e.g. Maynard Smith and Parker, 1976; Hammerstein and Parker, 1982), involving cases with both correlated and uncorrelated asymmetries.

The ESS was defined for a single population only and the stability condition of the original definition cannot be easily extended for bimatrix games. This is because bimatrix games assume that individuals of one species interact with individuals of the other species only, so there is no frequency
dependent mechanism that could prevent mutants of one population from invading residents of that population at the two species NE. In fact, it was shown (Selten, 1980) that requiring the stability condition of the ESS definition to hold in bimatrix games restricts the ESSs to strict NEs, i.e., to pairs of pure strategies. Two key assumptions behind Selten’s Theorem are that the probability that an individual occupies a given role is not affected by the strategies that it employs, and that payoffs within a given role are linear, as in matrix games. If either of these assumptions are violated then mixed strategy ESSs can result (see e.g. Webb et al, 1999; Broom and Rychtar, 2013).

There are interior NE in bimatrix games that deserve to be called “stable”, albeit in a weaker sense than was used in the (single-species) ESS definition. For example some of the NEs are stable with respect to some evolutionary dynamics (e.g., with respect to the replicator dynamics, or the best response dynamics). A static concept that captures such stability that proved useful for bimatrix games is the Nash–Pareto equilibrium (Hofbauer and Sigmund, 1998). The Nash-Pareto equilibrium is a NE which satisfies an additional condition that says that it is impossible for both players to increase their fitness by deviating from this equilibrium. For two-species games that cannot be described by a bimatrix (e.g., see Section 1.7.4), this concept of 2-species evolutionary stability was generalized by Cressman (2003) (see Volume I, Chapter Evolutionary games) who defined a 2-species ESS \((p^*, q^*)\) as a NE such that, if the population distributions of the two species are slightly perturbed then an individual in at least one species does better by playing its ESS strategy than by playing the slightly perturbed strategy of this species. We illustrate these concepts in the next section.

### 1.6.1 The Battle of the Sexes

A classical example of an asymmetric game is the *Battle of the Sexes* (Dawkins, 1976), where a population contains females with two strategies, coy and fast, and males with two strategies, faithful and philanderer. A coy female needs a period of courtship, whereas a fast female will mate with a male as soon as they meet. Faithful males are willing to engage in long courtships, and after mating will care for the offspring. A philanderer male will not engage in courtship, and so cannot mate with a coy female, and also leaves immediately after mating with a fast female.

Clearly in this case, any particular individual always occupies a given role (i.e. male or female) and cannot switch roles as is the case in the Owner-Intruder game in Section 1.6.2 below. Thus, males and females each have their own payoff matrix which are often represented as a bimatrix. The payoff bimatrix for the Battle of the Sexes is:
Here $B$ is the fitness gained by having an offspring, $C_R$ is the (potentially shared) cost of raising the offspring and $C_C$ is the cost of engaging in a courtship. All three of these terms are clearly positive. The above bimatrix is written in the form $(A_1, A_2^T)$, where matrix $A_1$ is the payoff matrix for males (player 1) and matrix $A_2$ is the payoff matrix for females (player 2), respectively.

For such games, to define a two-species NE, we study position of the two equal payoff lines, one for each sex. The equal payoff line for males (see the horizontal dashed line in Figure 1.1) is defined to be those $(p, q) \in \Delta_2 \times \Delta_2$ for which the payoff when playing Faithful equals the payoff when playing the Philanderer strategy, i.e.,

$$(1, 0)A_1q^T = (0, 1)A_1q^T$$

which yields $q_1 = \frac{C_R}{2(B - C_C)}$. Similarly, along the equal payoff line for females (see the vertical dashed line in Figure 1.1) the payoff when playing strategy Coy must equal the payoff when playing strategy Fast, i.e.,

$$(1, 0)A_2p^T = (0, 1)A_2p^T$$

If the two equal payoff lines do not intersect in the unit square, no completely mixed strategy (both for males and females) is a NE (Figure 1.1 A,B). In fact, there is a unique ESS (Philanderer, Coy), i.e., with no mating, for sufficiently small $B$ (Figure 1.1A), $B < \min(C_R/2 + C_C, C_R)$ (clearly not appropriate for a real population), a unique ESS (Philanderer, Fast) for sufficiently high $B$ (Figure 1.1B), when $B > C_R$. For intermediate $B$ satisfying $C_R/2 + C_C < B < C_R$ there is a 2-species weak ESS

$$p = \left(\frac{B - C_R}{B - C_C - C_R}, \frac{C_C}{C_C + C_R - B}\right), \quad q = \left(\frac{C_R}{2(B - C_C)}, 1 - \frac{C_R}{2(B - C_C)}\right),$$

where at least the fitness of one species increases toward this equilibrium, except when $p_1 = \frac{B - C_R}{B - C_C - C_R}$ or $q_1 = \frac{C_R}{2(B - C_C)}$ (Figure 1.1C). In all three cases of Figure 1.1, the NE is a Nash-Pareto pair, because it is impossible for both players to simultaneously deviate from the Nash equilibrium and increase their payoffs. In panels A and B both arrows point in direction of the NE. In panel C at least one arrow is pointing from the NE $(p, q)$ if both players deviate from that equilibrium. However, this interior NE is not a 2-species ESS since, when only one player (e.g., player one) deviates, sometimes no arrow points in the direction of $(p, q)$. This happens on the equal payoff lines (dashed lines). For example, let us consider points on the vertical dashed line above the NE. Here vertical arrows are zero vectors and horizontal arrows
point away from \( p \). Excluding points on the vertical and horizontal line from the definition of a 2-species ESS leads to a 2-species weak ESS.

Fig. 1.1 The ESS for the Battle of the Sexes game. Panel A assumes small \( B \) and the only ESS is \((p_1, q_1) = (0,1)\) = (Philanderer, Coy). Panel B assumes large \( B \) and the only ESS is \((p_1, q_1) = (0,0)\) = (Philanderer, Fast). For intermediate values of \( B \) (panel C) there is an interior NE. The dashed lines are the two equal payoff lines for males (horizontal line) and females (vertical line). The direction in which the male and female payoffs increase are shown by arrows (e.g., a horizontal arrow to the right means the first strategy (Faithful) has the higher payoff for males whereas a downward arrow means the second strategy (Fast) has the higher payoff for females). We observe that in panel C these arrows are such that at least the payoff of one player increases toward the Nash–Pareto pair, with the exception of the points that lie on the equal payoff lines. This qualifies the interior NE as a 2-species weak ESS.

1.6.2 The owner intruder game

The owner intruder game is an extension of the Hawk-Dove game, where player 1 (the owner) and player 2 (the intruder) have distinct roles (i.e., they cannot be interchanged as is the case of symmetric games). In particular an individual can play either of Hawk or Dove in either of the two roles. This leads to the bimatrix representation of the Hawk-Dove game below, which cannot be collapsed down to the single \( 2 \times 2 \) matrix from Section 1.2, because the strategy that an individual plays may be conditional upon the role that it occupies (in Section 1.2 there are no such distinct roles). The bimatrix of the game is

\[
\begin{array}{c|cc}
\text{Owner} & \text{Hawk} & \text{Dove} \\
\hline
\text{Hawk} & \left( \frac{V-C}{2}, \frac{V-C}{2} \right) & (V,0) \\
\text{Dove} & (0,V) & (V/2,V/2)
\end{array}
\]

Provided we assume that each individual has the same chance to be an owner or an intruder, the game can be symmetrized with the payoffs to the symmetrized game given in the following payoff matrix,
where

\[
\begin{align*}
\text{Hawk} & : (V - C)/2 & V & (3V - C)/4 & (3V - C)/4 \\
\text{Dove} & : 0 & V/2 & V/4 & V/4 \\
\text{Bourgeois} & : (V - C)/4 & 3V/4 & V/2 & (2V - C)/4 \\
\text{Anti-Bourgeois} & : (V - C)/4 & 3V/4 & (2V - C)/4 & V/2
\end{align*}
\]

Hawk – play Hawk when both owner and intruder,
Dove – play Dove when both owner and intruder,
Bourgeois – play Hawk when owner and Dove when intruder,
Anti-Bourgeois – play Dove when owner and Hawk when intruder.

It is straightforward to show that if \( V \geq C \) then Hawk is the unique ESS (Figure 1.2A), and that if \( V < C \) then Bourgeois and Anti-Bourgeois (alternatively called Marauder) are the only ESSs (Figure 1.2B). As we see, there are only pure strategy solutions, as opposed to the case of the Hawk Dove game, which had a mixed ESS \( V/C \) for \( V < C \), because the interior NE in Figure 1.2B is not a Nash–Pareto pair as in the upper left and lower right regions both arrows are pointing away from the NE. Thus, if both players simultaneously deviate from the Nash equilibrium their payoffs increase. Consequently, this NE is not a 2-species (weak) ESS.

Fig. 1.2 The ESS for the owner intruder game. Panel A assumes \( V > C \) and the only ESS is to be Hawk at both roles (i.e., \((p_1, q_1) = (1, 1) = (\text{Hawk, Hawk})\)). If \( V < C \) (Panel B), there are two boundary ESSs (black dots) corresponding to the Bourgeois \((p_1, q_1) = (1, 0) = (\text{Hawk, Dove})\) and Anti-Bourgeois \((p_1, q_1) = (0, 1) = (\text{Dove, Hawk})\) strategies. The direction in which the owner and intruder payoffs increase are shown by arrows (e.g., a horizontal arrow to the right means the Hawk strategy has the higher payoff for owner whereas a downward arrow means the Dove strategy has the higher payoff for intruder). The interior NE (the light gray dot at the intersection of the two equal payoff lines) is not a 2-species ESS as there are regions (the upper-left and lower-right corners) where both arrows point in directions away from this point.

Important recent work on this model and its ramifications for the part that respecting ownership plays has been carried out by Mesterton Gibbons,
Sherratt and co-workers (see Mesterton-Gibbons and Sherratt, 2014; Sherratt and Mesterton-Gibbons, 2015). In particular, why in real populations is the Bourgeois respect for ownership strategy so common, and the Anti-Bourgeois strategy so rare? One explanation offered by Maynard Smith (1982) was “infinite regress”. In this argument, immediately after a contest the winner becomes the owner of the territory and the loser becomes a potential intruder which could immediately rechallenge the individual that has just displaced it. In an Anti-Bourgeois population, this would result in the new owner conceding and the new intruder (the previous owner) once again being the owner; but then the displaced owner could immediately rechallenge, and the process would continue indefinitely. It is shown in Mesterton-Gibbons and Sherratt (2014) that under certain circumstances, but not always, this allows Bourgeois to be the unique ESS. Sherratt and Mesterton-Gibbons (2015) discuss many issues, such as uncertainty of ownership, asymmetry of resource value, continuous contest investment (as in the war of attrition) and potential signalling of intentions (what they call “secret handshakes”, similar to some of the signals we discuss in Section 1.10) in detail. There are many reasons that can make evolution of Anti-Bourgeois unlikely, and it is probably a combination of these that make it so rare.

1.6.3 Bimatrix replicator dynamics

The single species replicator dynamics such as those for the Hawk-Dove game (Section 1.2.1) can be extended to two roles as follows (Hofbauer and Sigmund, 1998). Note that here this is interpreted as two completely separate populations, i.e., any individual can only ever occupy one of the roles, and its offspring occupy that same role. If $A = (a_{ij})_{i=1,...,n}$ and $B = (b_{ij})_{j=1,...,m}$ are the payoff matrices to an individual in role 1 and role two, respectively, the corresponding replicator dynamics are

$$\frac{dp_{1i}(t)}{dt} = p_{1i}((Ap_{2}^T)_i - p_1 Ap_2^T) \quad i = 1,...,n;$$

$$\frac{dp_{2j}(t)}{dt} = p_{2j}((Bp_{1}^T)_j - p_2 Bp_1^T) \quad j = 1,...,m;$$

where $p_1 \in \Delta_n$ and $p_2 \in \Delta_m$ are the population mixtures of individuals in role 1 and 2, respectively. For example for the two role, two strategy game, where without loss of generality we can set $a_{11} = a_{22} = b_{11} = b_{22} = 0$ (since as for matrix games, adding a constant to all of the payoffs an individual gets against a given strategy does not affect the NEs/ESSs), we obtain
\[
\frac{dx}{dt} = x(1-x)(a_{12} - (a_{12} + a_{21})y),
\]
\[
\frac{dy}{dt} = y(1-y)(b_{12} - (b_{12} + b_{21})x),
\]  

where \(x\) is the frequency of the first strategy players in the role 1 population, and \(y\) is the corresponding frequency for role 2. Hofbauer and Sigmund (1998) show that orbits converge to the boundary in all cases except if \(a_{12}a_{21} > 0\), \(b_{12}b_{21} > 0\) and \(a_{12}b_{12} < 0\), which yield closed periodic orbits around the internal equilibrium. Replicator dynamics for the Battle of the Sexes and the owner intruder game are shown in Figure 1.3.

Note there are some problems with the interpretation of the dynamics of two populations in this way, related to the assumption of exponential growth of populations, since the above dynamics effectively assume that the relative size of the two populations remains constant (Argasinski, 2006).

![Fig. 1.3 Bimatrix replicator dynamics (1.11) for the Battle of the Sexes game (A–C) and the owner intruder game (D, E), respectively. Panels A–C correspond to panels given in Figure 1.1 and panels D and E correspond to those of Figure 1.2. This figure shows that trajectories of the bimatrix replicator dynamics converge to 2-species ESS as defined in Sections 1.6.1 and 1.6.2. In particular, the interior NE in panel E is not a 2-species ESS and it is an unstable equilibrium for the bimatrix replicator dynamics. In panel C the interior NE is 2-species weak ESS and it is (neutrally) stable for the bimatrix replicator dynamics.](image-url)
1.7 The Habitat Selection game

Fretwell and Lucas (1969) introduced the Ideal Free Distribution (IFD) to describe a distribution of animals in a heterogeneous environment consisting of discrete patches $i = 1, \ldots, n$. The IFD assumes that animals are free to move between several patches, travel is cost-free, each individual knows perfectly the quality of all patches and all individuals have the same competitive abilities. Assuming that these patches differ in their basic quality $B_i$ (i.e., their quality when unoccupied), the IFD model predicts that the best patch will always be occupied.

Let us assume that patches are arranged in descending order ($B_1 > \cdots > B_n > 0$) and $m_i$ is the animal abundance in patch $i$. Let $p_i = m_i / (m_1 + \cdots + m_n)$ be the proportion of animals in patch $i$, so that $p = (p_1, \ldots, p_n)$ describes the spatial distribution of the population. For a monomorphic population, $p_i$ also specifies the individual strategy as the proportion of the lifetime an average animal spends in patch $i$. We assume that the payoff in each patch, $V_i(p_i)$, is a decreasing function of animal abundance in that patch, i.e., the patch payoffs are negatively density dependent. Then, fitness of a mutant with strategy $\tilde{p} = (\tilde{p}_1, \ldots, \tilde{p}_n)$ in the resident monomorphic population with distribution $p = (p_1, \ldots, p_n)$ is

$$E(\tilde{p}, p) = \sum_{i=1}^{n} \tilde{p}_i V_i(p_i).$$

However, we do not need to make the assumption that the population is monomorphic, because what really matters in calculating $E(\tilde{p}, p)$ above is the animal distribution $p$. If the population is not monomorphic this distribution can be different from strategies animals use and we call it the population mean strategy. Thus, in the Habitat Selection game individuals do not enter pairwise conflicts, but they play against the population mean strategy (referred to as a “playing the field”, or “population” game).

Fretwell and Lucas (1969) introduced the concept of the Ideal Free Distribution which is a population distribution $p = (p_1, \ldots, p_n)$ that satisfies two conditions:

1. There exists a number $1 \leq k \leq n$ such that $p_1 > 0, \ldots, p_k > 0$ and $p_{k+1} = \cdots = p_n = 0$.
2. $V_1(p_1) = \cdots = V_k(p_k) = V^*$ and $V^* \geq V_i(p_i)$ for $i = k + 1, \ldots, n$.

They proved that provided patch payoffs are negatively density dependent (i.e., decreasing functions of the number of individuals in a patch), then there exists a unique IFD which Cressman and Krivan (2006) later showed is an ESS. In the next section we will discuss two commonly used types of patch payoff functions.
1.7.1 Parker’s matching principle

Parker (1978) considered the case where resource input rates \( r_i \), \( i = 1, \ldots, n \) are constant and resources are consumed immediately when they enter the patch and so there is no standing crop. This leads to a particularly simple definition of animal patch payoffs as the ratio of the resource input rate divided by the number of individuals there, i.e.,

\[
V_i = \frac{r_i}{m_i} = \frac{r_i}{p_iM}
\]  

(1.12)

where \( M = m_1 + \cdots + m_n \) is the overall population abundance. The matching principle then says that animals distribute themselves so that their abundance in each patch is proportional to the rate with which resources arrive into the patch, \( p_i/p_j = r_i/r_j \). This is nothing other than the IFD for payoff functions (1.12). It is interesting to notice that all patches will be occupied independently of the total population abundance. Indeed, as the consumer density in the \( i \)-th patch decreases, payoff \( r_i/(p_iM) \) increases, which attracts some animals, and there cannot be unoccupied patches. There is an important difference between this (nonlinear) payoff function (1.12) and the linear payoff function that we consider in the following section (1.13), because as the local population abundance in a patch decreases, then (1.12) tends to infinity, but (1.13) tends to \( r_i \). This means that in the first case there cannot be unoccupied patches (irrespective of their basic patch quality \( r_i \)) because the payoffs in occupied patches are finite, but the payoff in unoccupied patches would be infinite (provided all \( r_i > 0 \)). This argument does not apply in the case of the logistic payoff (1.13). This concept successfully predicts the distribution of house flies that arrive at a cow pat where they immediately mate (Parker, 1978, 1984; Blanckenhorn et al, 2000), or of fish that are fed at two feeders in a tank (Milinski, 1979, 1988; Berec et al, 2006).

In the next section we consider the situation where resources are not consumed immediately upon entering the system.

1.7.2 Patch payoffs are linear

Here we consider two patches only and we assume that the payoff in habitat \( i (= 1, 2) \) is a linearly decreasing function of population abundance

\[
V_i = r_i \left( 1 - \frac{m_i}{K_i} \right) = r_i \left( 1 - \frac{p_iM}{K_i} \right)
\]  

(1.13)

where \( m_i \) is the population density in habitat \( i \), \( r_i \) is the intrinsic per capita population growth rate in habitat \( i \), and \( K_i \) is its carrying capacity. The total population size in the two habitat environment is denoted by \( M (= m_1 + m_2) \).
and the proportion of the population in habitat \( i \) is \( p_i = m_i/M \). Payoff (1.13) is often used in population dynamics where it describes the logistic population growth.

Let us consider an individual which spends proportion \( \tilde{p}_1 \) of its lifetime in habitat 1 and \( \tilde{p}_2 \) in habitat 2. Provided total population density is fixed at \( M \), then its fitness in the population with mean strategy \( \mathbf{p} = (p_1, p_2) \) is

\[
E(\mathbf{p}, \mathbf{p}) = \tilde{p}_1 V_1(p_1) + \tilde{p}_2 V_2(p_2) = \tilde{p} \mathbf{U} \mathbf{p}^T,
\]

where

\[
U = \begin{pmatrix}
    r_1(1 - \frac{M}{K_1}) & r_1 \\
    r_2(1 - \frac{M}{K_2}) & r_2
\end{pmatrix}
\]

is the payoff matrix with two strategies, where strategy \( i \) represents staying in patch \( i \) \( (i = 1, 2) \). This shows that the Habitat Selection game with a linear payoff can be written for a fixed population size as a matrix game. If the per capita intrinsic population growth rate in habitat 1 is higher than that in habitat 2 \( (r_1 > r_2) \) the IFD is (Křivan and Sirot, 2002)

\[
p_1 = \begin{cases} 
1 & \text{if } M < K_1 \frac{r_1 - r_2}{r_1} \\
\frac{r_2 K_1}{r_2 K_2} + \frac{K_1 K_2 (r_1 - r_2)}{(r_2 K_1 + r_1 K_2) M} & \text{otherwise.}
\end{cases} 
\]

(1.14)

When the total population abundance is low, the payoff in habitat 1 is higher than the payoff in habitat 2 for all possible population distributions because the competition in patch 1 is low due to low population densities. For higher population abundances, neither of the two habitats is always better than the other, and under the IFD payoffs in both habitats must be the same \( (V_1(p_1) = V_2(p_2)) \). Once again, it is important to emphasize here that the IFD concept is different from maximization of the mean animal fitness \( \overline{W}(\mathbf{p}, \mathbf{p}) = p_1 V_1(p_1) + p_2 V_2(p_2) \)

which would lead to

\[
p_1 = \begin{cases} 
1 & \text{if } M < K_1 \frac{r_1 - r_2}{2r_1} \\
\frac{r_2 K_1}{r_1 K_2 + r_2 K_1} + \frac{K_1 K_2 (r_1 - r_2)}{2(r_1 K_2 + r_2 K_1) M} & \text{otherwise.}
\end{cases} 
\]

(1.15)

The two expressions (1.14) and (1.15) are the same if and only if \( r_1 = r_2 \). Interestingly, by comparing (1.14) and (1.15), we see that maximizing mean fitness leads to fewer animals than the IFD in the patch with higher basic quality \( r_i \) (i.e., in patch 1).
1.7.3 Some extensions of the Habitat Selection game

The Habitat Selection game as described makes several assumptions that were relaxed in the literature. One assumption is that patch payoffs are decreasing functions of population abundance. This assumption is important because it guarantees that a unique IFD exists. However, patch payoffs can also be increasing functions of population abundance. In particular, at low population densities payoffs can increase as more individuals enter a patch and competition is initially weak. For example, more individuals in a patch can increase the probability of finding a mate. This is called the Allee effect. The IFD for the Allee effect has been studied in the literature (Fretwell and Lucas, 1969; Morris, 2002; Krivan, 2014; Cressman and Tran, 2015). It has been shown that for hump shaped patch payoffs, up to three IFDs can exist for a given overall population abundance. At very low overall population abundances, only the most profitable patch will be occupied. At intermediate population densities, there are two IFDs corresponding to pure strategies where all individuals occupy patch 1 only, or patch 2 only. As population abundance increases, competition becomes more severe and an interior IFD appears exactly as in the case of negative density dependent payoff functions. At high overall population abundances only the interior IFD exists due to strong competition among individuals. It is interesting to note that as the population numbers change, there can be sudden (discontinuous) changes in the population distribution. Such erratic changes in the distribution of deer mice were observed and analyzed by Morris (2002).

Another complication that leads to multiple IFDs is the cost of dispersal. Let us consider a positive migration cost $c$ between two patches. An individual currently in patch 1 will migrate to patch 2 only if the payoff there is such that $V_2(p_2) - c \geq V_1(p_1)$. Similarly, an individual currently in patch 2 will migrate to patch 1 only if its payoff does not decrease by doing so, i.e., $V_1(p_1) - c \geq V_2(p_2)$. Thus, all distributions $(p_1, p_2)$ that satisfy these two inequalities form the set of IFDs (Mariani et al, 2016).

The Habitat selection game was also extended to situations where individuals perceive space as a continuum (e.g., Cosner, 2005; Cantrell et al, 2007, 2012). The movement by diffusion is then combined, or replaced, by a movement along the gradient of animal fitness.

1.7.4 Habitat selection for two species

Instead of a single species we now consider two species with population densities $M$ and $N$ dispersing between two patches. We assume that individuals of these species compete in each patch both intra- as well as inter-specifically. Following our single-species Habitat Selection game, we assume that individual payoffs are linear functions of species distribution (Krivan and Sirot, ...
\[ V_i(p, q) = r_i \left( 1 - \frac{p_i M}{K_i} - \frac{\alpha_i q_i N}{K_i} \right), \]
\[ W_i(p, q) = s_i \left( 1 - \frac{q_i N}{L_i} - \frac{\beta_i p_i M}{L_i} \right), \]

where \( p = (p_1, p_2) \) denotes the distribution of species 1 and \( q = (q_1, q_2) \) the distribution of species 2. Here, positive parameters \( \alpha_i \) (respectively \( \beta_i \)), are interspecific competition coefficients, \( r_i \) (respectively \( s_i \)) are the intrinsic per capita population growth rates and \( K_i \) (respectively \( L_i \)) are the environmental carrying capacities. The two-species Habitat Selection game cannot be represented in a bimatrix form (to represent it in a matrix form we would need four matrices), because the payoff in patch \( i \) for a given species depends not only on the distribution (strategy) of its competitors, but also on the distribution of its own conspecifics. The equal payoff line for species one (two) are those \((p, q) \in \Delta_2 \times \Delta_2 \) for which \( V_1(p, q) = V_2(p, q) \) (\( W_1(p, q) = W_2(p, q) \)).

Since payoffs are linear functions these are lines in the coordinates \( p_1 \) and \( q_1 \), but as opposed to the case of bimatrix games in Section 1.6.1 and 1.6.2 they are neither horizontal, nor vertical. If they do not intersect in the unit square, the two species cannot coexist in both patches at a NE. The most interesting case is when the two equal payoff lines intersect inside the unit square. Krivan et al (2008) showed that the interior intersection is the 2-species ESS provided
\[
\begin{align*}
  r_1 s_1 K_2 L_2 (1 - \alpha_1 \beta_1) &+ r_1 s_2 K_3 L_1 (1 - \alpha_1 \beta_2) + \\
  r_2 s_1 K_1 L_2 (1 - \alpha_2 \beta_1) &+ r_2 s_2 K_1 L_1 (1 - \alpha_2 \beta_2) > 0.
\end{align*}
\]

Geometrically, this condition states that the equal payoff line for species one has a more negative slope than that for species two. This allows us to extend the concept of the single-species Habitat Selection game to two species that compete in two patches. In this case the 2-species IFD is defined as a 2-species ESS. We remark that the best response dynamics do converge to such two-species IFD (Krivan et al, 2008).

One of the predictions of the Habitat selection game for two species is that as competition gets stronger, the two species will spatially segregate (e.g., Morris, 1999; Krivan and Sirot, 2002). Such spatial segregation was observed in experiments with two bacterial strains in a microhabitat system with nutrient poor and nutrient rich patches (Lambert et al, 2011).
1.8 Dispersal and evolution of dispersal

Organisms often move from one habitat to another, which is referred to as dispersal. We focus here on dispersal and its relation to the IFD discussed in the previous section.

Here we consider \( n \) habitat patches and a population of individuals that disperse between them. In what follows we will assume that the patches are either adjacent (in particular when there are just two patches), or that the travel time between them is negligible when compared to the time individuals spend in these patches. There are two basic questions:

1. When is dispersal an adaptive strategy, i.e., when does individual fitness increase for dispersing animals compared to those who are sedentary?
2. Where should individuals disperse?

To describe changes in population densities we will consider demographic population growth in each patch and dispersal between patches. Dispersal is described by the propensity of individuals to disperse \( (\delta \geq 0) \) and by a dispersal matrix \( D \). The entries of this matrix \( (D_{ij}) \) describe the transition probabilities that an individual currently in patch \( j \) moves to patch \( i \) per unit of time. We remark that \( D_{ii} \) is the probability of staying in patch \( i \). Per capita population growth rates in patches are given by \( f_i \) (for example, \( f_i \) can be the logistic growth rate \( V_i \) (1.13) in Section 1.7.2). The changes in population numbers are then described by population–dispersal dynamics

\[
\frac{dm_i}{dt} = m_i f_i(m_i) + \delta \sum_{j=1}^{n} (D_{ij}(m)m_j - D_{ji}(m)m_i) \quad \text{for} \quad i = 1, \ldots, n \quad (1.16)
\]

where \( m = (m_1, \ldots, m_n) \) is the vector of population densities in \( n \) patches. Thus, the first term in the above summation describes immigration to patch \( i \) from other patches and the second term describes emigration from patch \( i \) to other patches. In addition, we assume that \( D \) is irreducible, i.e., there are no isolated patches.

The case that corresponds to the passive diffusion between patches assumes that entries of the dispersal matrix are constant and the matrix is symmetric. It was shown (Takeuchi, 1996) that when functions \( f_i \) are decreasing with \( f_i(0) > 0 \) and \( f_i(K_i) = 0 \) for some \( K_i > 0 \) then model (1.16) has an interior equilibrium which is globally asymptotically stable. However this does not answer the question of whether such an equilibrium is evolutionarily stable, i.e., whether it is resistant to invasion of mutants with the same traits (parameters) as the resident population, but different propensity to disperse \( \delta \).

The answer to this question depends on the entries of the dispersal matrix. An interior population distribution \( m^* = (m_1^*, \ldots, m_n^*) \) will be the IFD provided patch payoffs in all patches are the same, i.e., \( f_1(m_1^*) = \cdots = f_n(m_n^*) \).

This implies that at the population equilibrium there is no net dispersal, i.e.,
\[ \delta \sum_{j=1}^{n} (D_{ij}m_j^* - D_{ji}m_i^*) = 0. \]

There are two possibilities. Either

\[ \sum_{j=1}^{n} (D_{ij}m_j^* - D_{ji}m_i^*) = 0, \quad (1.17) \]

or \( \delta = 0 \). The pattern of equalized immigration and emigration satisfying (1.17) is called “balanced dispersal” (McPeek and Holt, 1992; Doncaster et al, 1997; Holt and Barfield, 2001). Under balanced dispersal, there is an inverse relation between local population size and its dispersal rate. In other words, individuals at good sites are less likely to disperse than those from poor sites. When dispersal is unbalanced, Hastings (1983) showed that mutants with lower propensity to disperse will outcompete the residents and no dispersal (\( \delta = 0 \)) is the only evolutionarily stable strategy.

However, dispersal can be favored even when it is not balanced. Hamilton and May (1977) showed that unconditional and costly dispersal among very many patches can be promoted because it reduces competition between relatives. Their model was generalized by Comins et al (1980) who assumed that because of stochastic effects a proportion \( e \) of patches can become empty at any time step. A proportion \( p \) of migrants survives migration and redistributes at random (assuming the Poisson distribution) among the patches. These authors derived analytically the evolutionarily stable dispersal strategy that is given by a complicated implicit formula (see formula (3) in Comins et al, 1980). As population abundance increases the evolutionarily stable dispersal rate converges to a simpler formula

\[ \delta = \frac{e}{1 - p(1 - e)}. \]

Here the advantage of dispersal results from the possibility of colonizing an extinct patch.

Evolution of mobility in predator-prey systems was also studied by Xu et al (2014). These authors showed how interaction strength between mobile vs. sessile prey and predators influences the evolution of dispersal.

### 1.9 Foraging games

Foraging games describe interactions between prey, their predators, or both. These games assume that either predator or prey behave in order to maximize their fitness. Typically, the prey strategy is to avoid predators while predators try to track their prey. Several models that focus on various aspects of
predator-prey interactions were developed in the literature (e.g., Brown and Vincent, 1987; Brown et al, 1999, 2001; Vincent and Brown, 2005).

An important component of predation is the functional response defined as the per predator rate of prey consumption (Holling, 1959). It also serves as a basis for models of optimal foraging (Stephens and Krebs, 1986) that aim to predict diet selection of predators in environments with multiple prey types. In this section we start with a model of optimal foraging and we show how it can be derived using extensive form games (see Volume I, Chapter Evolutionary games). As an example of a predator-prey game we then discuss predator-prey distribution in a two-patch environment.

**Fig. 1.4** The decision tree for two prey types. The first level gives the prey encounter distribution. The second level gives the predator activity distribution. The final row of the diagram gives the probability of each predator activity event and so sums to 1. If prey 1 is the more profitable type, the edge in the decision tree corresponding to not attacking this type of prey is never followed at optimal foraging (indicated by the dashed edge in the tree).

### 1.9.1 Optimal foraging as an agent normal form of an extensive game

Often it is assumed that a predator’s fitness is proportional to its prey intake rate and the functional response serves as a proxy of fitness. In the case of two or more prey types, the multi-prey functional response is the basis of the diet choice model (Charnov, 1976; Stephens and Krebs, 1986) that predicts the predator’s optimal diet as a function of prey densities in the environment.
Here we show how functional responses can be derived using decision trees of games given in extensive form (Cressman, 2003; Cressman et al., 2014) (see also Broom et al., 2004, for an example of where this methodology was used in a model of food stealing). Let us consider a decision tree in Figure 1.4 describing a single predator feeding on two prey types. This decision tree assumes that a searching predator meets prey type 1 with probability \( p_1 \) and prey type 2 with probability \( p_2 \) during the search time \( \tau \). For simplicity we will assume that \( p_1 + p_2 = 1 \).

Upon an encounter with a prey individual, the predator decides whether to attack the prey (prey type 1 with probability \( q_1 \) and prey type 2 with probability \( q_2 \)) or not. When a prey individual is captured, the energy that the predator receives is denoted by \( e_1 \) or \( e_2 \). The predator’s cost is measured by the time lost. This time consists of the search time \( \tau \) and the time needed to handle the prey (\( h_1 \) for prey type 1 and \( h_2 \) for prey type 2).

Calculation of functional responses is based on renewal theory which proves that the long term intake rate of a given prey type can be calculated as the mean energy intake during one renewal cycle divided by the mean duration of the renewal cycle (Stephens and Krebs, 1986; Houston and McNamara, 1999). A single renewal cycle is given by a predator passing through the decision tree in Figure 1.4. Since type \( i \) prey are only killed when the path denoted by \( p_i \) and then \( q_i \) is followed, the functional response to prey \( i = 1, 2 \) is:

\[
 f_i(q_1, q_2) = \frac{p_i q_i}{\tau + p_i q_1 h_1 + p_2 q_2 h_2}.
\]

When \( x_i \) denotes density of prey type \( i \) in the environment and the predator meets prey at random, \( p_i = x_i/x \), where \( x = x_1 + x_2 \). Setting \( \lambda = 1/(\tau_x) \) leads to

\[
 f_i(q_1, q_2) = \frac{\lambda x_i q_i}{1 + \lambda x_1 q_1 h_1 + \lambda x_2 q_2 h_2}.
\]

These are the functional responses assumed in standard two prey type models. The predator’s rate of energy gain is given by

\[
 E(q_1, q_2) = e_1 f_1(q_1, q_2) + e_2 f_2(q_1, q_2) = \frac{e_1 p_1 q_1 + e_2 p_2 q_2}{\tau + p_1 q_1 h_1 + p_2 q_2 h_2}, \quad (1.18)
\]

This is the proxy of the predator’s fitness which is maximized over the predator’s diet \( (q_1, q_2) \), \( 0 \leq q_i \leq 1, i = 1, 2 \).

Here, using the agent normal form of extensive form game theory (Cressman, 2003), we show an alternative, game theoretical, approach to find the optimal foraging strategy. This method assigns a separate player (called an agent) to each decision node (here 1 or 2). The possible decisions at this node become the agent’s strategies and its payoff is given by the total energy intake rate of the predator it represents. Thus, all of the virtual agents have the
same common payoff. The optimal foraging strategy of the single predator is then a solution to this game. In our example, player 1 corresponds to decision node 1 with strategy set \( \Delta_1 = \{ q_1 | 0 \leq q_1 \leq 1 \} \) and player 2 to node 2 with strategy set \( \Delta_2 = \{ q_2 | 0 \leq q_2 \leq 1 \} \). Their common payoff \( E(q_1, q_2) \) is given by (1.18) and we seek the NE of the two-player game. Assuming that prey type 1 is the more profitable for the predator, its energy content per unit handling time is higher than the profitability of the second prey type (i.e., \( e_1/h_1 > e_2/h_2 \) we get \( E(1, q_2) > E(q_1, q_2) \) for all \( 0 \leq q_1 < 1 \) and \( 0 \leq q_2 \leq 1 \). Thus, at any NE, player 1 must play \( q_1 = 1 \). The NE strategy of player 2 is then any best response to \( q_1 = 1 \) (i.e., any \( q_2 \) that satisfies \( E(1, q_2') \leq E(1, q_2) \) for all \( 0 \leq q_2' \leq 1 \) which yields

\[
q_2 = \begin{cases} 
0 & \text{if } p_1 > p_1^* \\
1 & \text{if } p_1 < p_1^* \\
[0, 1] & \text{if } p_1 = p_1^*,
\end{cases}
\]

where

\[
p_1^* = \frac{e_2 \tau_s}{e_1 h_2 - e_2 h_1}.
\]

This NE coincides with the optimal strategy derived by maximization of (1.18). It makes quite striking predictions. While the more profitable prey type is always included in the predator’s diet, inclusion of the less profitable prey type is independent of its own density and depends on the more profitable prey type density only. This prediction was experimentally tested with great tits (e.g., Krebs et al, 1977; Berec et al, 2003). That the Nash equilibrium coincides with the optimal foraging strategy (i.e., with the maximum of \( E \)) in this model is not a coincidence. Cressman et al (2014) proved that this is so for all foraging games with a 2-level decision tree. For decision trees with more levels, they showed that the optimal foraging strategy is always a NE of the corresponding agent normal form game and that other, non-optimal, NE may also appear.

### 1.9.2 A predator-prey foraging game

As an example we consider here a predator-prey foraging game between prey and predators in a two-patch environment. If \( x_i \) denotes the abundance of prey in patch \( i (= 1, 2) \), the total abundance of prey is \( x = x_1 + x_2 \) and, similarly, the total abundance of predators is \( y = y_1 + y_2 \). Let \( u = (u_1, u_2) \) be the distribution of prey and \( v = (v_1, v_2) \) be the distribution of predators. We neglect the travel time between patches so that \( u_1 + u_2 = v_1 + v_2 = 1 \) (i.e., each animal is either in patch 1 or patch 2). We assume that the prey population grows exponentially at each patch with the per capita population growth rate \( r_i \) and it is consumed by predators. The killing rate is given by
the functional response. For simplicity we neglect the handling time so that the functional response in patch $i$ is $f_i = \lambda_i x_i$, i.e., the per prey per predator killing rate is $\lambda_i$. The prey payoff in patch $i$ is given by the per capita prey population growth rate in that patch, i.e., $r_1 - \lambda_1 v_1 y$ as there are $v_1 y$ predators in patch $i$. The fitness of a prey individual is

$$V(u, v) = (r_1 - \lambda_1 v_1 y)u_1 + (r_2 - \lambda_2 v_2 y)u_2.$$  \hfill (1.21)

The predator payoff in patch $i$ is given by the per capita predator population growth rate $e_i u_i x - m_i$, where $e_i$ is a coefficient by which the energy gained by feeding on prey is transformed into new predators and $m_i$ is the per capita predator mortality rate in patch $i$. The fitness of a predator with strategy $v = (v_1, v_2)$ when the prey use strategy $u = (u_1, u_2)$ is

$$W(v, u) = (e_1 \lambda_1 u_1 x - m_1) v_1 + (e_2 \lambda_2 u_2 x - m_2) v_2.$$  \hfill (1.22)

This predator-prey game can be represented by the following payoff bimatrix

$$
\begin{array}{c|cc}
\text{Prey} & \text{Patch 1} & \text{Patch 2} \\
\hline
\text{Patch 1} & (r_1 - \lambda_1 y, e_1 \lambda_1 x - m_1) & (r_1, -m_2) \\
\text{Patch 2} & (r_2, -m_1) & (r_2 - \lambda_2 y, e_2 \lambda_2 x - m_2)
\end{array}
$$

That is, the rows in this bimatrix correspond to the prey strategy (the first row means the prey are in patch 1, the second row means prey are in patch 2) and similarly columns represent the predator strategy. The first of the two expressions in the entries of the bimatrix is the payoff for the prey, the second is payoff for the predators.

For example, we will assume that for prey patch 1 has a higher basic patch quality when compared to patch 2 (i.e., $r_1 \geq r_2$) while for predators patch 1 has a higher mortality rate ($m_1 > m_2$). The corresponding NE is (Křivan, 1997)

(a) $(u_1^*, v_1^*)$ if $x > \frac{m_1 - m_2}{e_1 \lambda_1}$, $y > \frac{r_1 - r_2}{\lambda_1}$,

(b) $(1, 1)$ if $x > \frac{m_1 - m_2}{e_1 \lambda_1}$, $y < \frac{r_1 - r_2}{\lambda_1}$,

(c) $(1, 0)$ if $x < \frac{m_1 - m_2}{e_1 \lambda_1}$,

where

$$(u_1^*, v_1^*) = \left( \frac{m_1 - m_2 + e_2 \lambda_2 x}{(e_1 \lambda_1 + e_2 \lambda_2) x}, \frac{r_1 - r_2 + \lambda_2 y}{(\lambda_1 + \lambda_2 y)} \right).$$

If prey abundance is low (case (c)), all prey will be in patch 1 while predators will stay in patch 2. Because the mortality rate for predators in patch 1 is higher than in patch 2 and prey abundance is low, patch 2 is a refuge for predators. If predator abundance is low and prey abundance is high (case (b)), both predators and prey will aggregate in patch 1. When the NE is
strict (cases (b) and (c) above) it is also the ESS because there is no alternative strategy with the same payoff. However, when the NE is mixed (case (a)), there exist alternative best replies to it. This mixed NE is the 2-species weak ESS. It is globally asymptotically stable for the continuous-time best response dynamics (Křivan et al, 2008) that model dispersal behavior whereby individuals move to the patch with the higher payoff. We remark that for some population densities \( x = \frac{m_1}{m_1 - m_2 e^{\lambda_1}} \) and \( y = \frac{r_1}{r_1 - r_2 e^{\lambda_1}} \) the NE is not uniquely defined, which is a general property of matrix games (the game is non-generic in this case).

1.10 Signalling games

Signalling between animals occurs in a number of contexts. This can be signals, often but not necessarily between conspecifics, warning of approaching predators. This situation can be game theoretic, as the signaller runs a potentially higher risk of being targeted by the predator. There are also cases of false signals being given when no predator is approaching to force food to be abandoned which can then be consumed by the signaller (Flower, 2011). Alternatively within a group of animals, each individual may need to decide how to divide their time between vigilance and feeding, where each individual benefits from the vigilance of others as well as itself, and this has been modelled game-theoretically (e.g., McNamara and Houston, 1992; Brown, 1999; Sirot, 2012).

Another situation occurs between relatives over items of food, for example a parent bird feeding its offspring. Young birds beg aggressively for food, and the parent must decide which to feed, if any (it can instead consume the item itself). The most well-known model of this situation is the Sir Philip Sidney game (Maynard Smith, 1991), and is a model of cost-free signalling (Bergstrom and Lachmann, 1998).

The classic example of a signalling game is between potential mates. Males of differing quality advertise this quality to females, often in a way that is costly, and the females choose who to mate with based upon the strength of the signal. Examples are the tail of the peacock or the elaborate bower created by bowerbirds. There is obviously a large incentive to cheat, and so how are such signals kept honest? A signal that is not at least partly correlated to quality would be meaningless, and so would eventually be ignored. The solution as developed by Zahavi (1975, 1977), the handicap principle, is that these costly signals are easier to bear by higher quality mates, and that evolution leads to a completely honest signal, where each quality level has a unique signal.
1.10.1 The game

The following signalling model is due to (Grafen, 1990a,b). Consider a population with a continuum of male quality types $q$ and a single type of female. Assume that a male of quality $q$ gives a signal $a = A(q)$ of this quality, where higher values of $a$ are more costly. It is assumed that there is both a minimum quality level $q_0 > 0$ (there may or may not be a maximum quality level) and a minimum signal level $a_0 \geq 0$ (which can be thought of as giving no signal). When a female receives a signal she allocates a quality level to the signal $P(a)$. We have a nonlinear asymmetric game with sequential decisions; in particular the nonlinearity makes this game considerably more complicated than asymmetric games such as the Battle of the Sexes of Section 1.6.1. The female pays a cost for misassessing a male of quality $q$ as being of quality $p$ of $D(q, p)$, which is positive for $p \neq q$, with $D(q, q) = 0$. Assuming that the probability density of males of quality $q$ is $g(q)$, the payoff to the female, which is simply minus the expected cost, is

$$- \int_{q_0}^{\infty} D(q, p) g(q) dq.$$

An honest signalling system with strategies $A^*$ and $P^*$ occurs if and only if $P^*(A^*(q)) = q$, for all $q$. We note that here the female never misassesses a male, and so pays zero cost. Clearly any alternative female assessment strategy would do worse. But how can we obtain stability against alternative (cheating) male strategies?

The fitness of a male of quality $q$, $W(a, p, q)$, depends upon his true quality, the quality assigned to him by the female and the cost of his signal. $W(a, p, q)$ will be increasing in $p$ and decreasing in $a$. For stability of the honest signal, we need that the incremental advantage of a higher level of signalling is greater for a high quality male than for a low quality one, so that

$$- \frac{\partial}{\partial a} W(a, p, q) \frac{\partial}{\partial p} W(a, p, q)$$

is strictly decreasing in $q$ (note that the ratio is negative, so minus this ratio is positive), i.e., the higher quality the male, the lower the ratio of the marginal cost to the marginal benefit for an increase in the level of advertising. This ensures that completely honest signalling cannot be invaded by cheating, since costs to cheats to copy the signals of better quality males would be explicitly higher than for the better quality males, who could always thus achieve a cost they were willing to pay that the lower quality cheats would not.

The following example male fitness function is given in Grafen (1990a) (the precise fitness function to the female does not affect the solution provided that correct assessment yields 0, and any misassessment yields a negative payoff)
with qualities in the range \( q_0 \leq q < 1 \) and signals of strength \( a \geq a_0 \), for some \( r > 0 \).

We can see that the function from (1.24) satisfies the above conditions on \( W(a, p, q) \). In particular consider the condition from expression (1.23)

\[
- \frac{\partial}{\partial a} W(a, p, q) = -p^r q^a \ln q, \quad \frac{\partial}{\partial p} W(a, p, q) = r p^{-1} q^a
\]

which are the increase in cost per unit increase in the signal level and the increase in the payoff per unit increase in the female’s perception (which in turn is directly caused by increases in signal level), respectively. The ratio from (1.23), which is proportional to the increase in cost per unit of benefit that this would yield, becomes \(-p \ln q / r\), takes a larger value for lower values of \( q \). Thus there is an honest signalling solution. This is shown in Grafen (1990a) to be given by

\[
A(q) = a_0 - r \ln \left( \frac{\ln(q)}{\ln(q_0)} \right), \quad P(a) = q_0^{\exp(-a-a_0)/r}.
\]

### 1.11 Conclusion

In this chapter we have covered some of the important evolutionary game models applied to biological situations. We should note that we have left out a number of important theoretical topics as well as areas of application. We briefly touch on a number of those below.

All of the games that we have considered involved either pairwise games, or playing the field games, where individuals effectively play against the whole population. In reality contests will sometimes involve groups of individuals. Such models were developed in Broom et al (1997), for a recent review see Gokhale and Traulsen (2014). In addition the populations were all both effectively infinite and well-mixed in the sense that for any direct contest involving individuals, each pair was equally likely to meet. In reality populations are finite and have (e.g., spatial) structure. The modelling of evolution in finite populations often uses the Moran process (Moran, 1958), but more recently games in finite populations have received significant attention (Nowak, 2006). These models have been extended to include population structure by considering evolution on graphs (Lieberman et al, 2005), and there has been an explosion of such model applications, especially to consider the evolution of cooperation. Another feature of realistic populations that we have ignored is the state of the individual. A hungry individual may behave differently to one that has recently eaten, and nesting behaviour may be different at the
start of the breeding season to later on. A theory of state-based models has
been developed in Houston and McNamara (1999).

In terms of applications we have focused on classical biological problems,
but game theory has also been applied to medical scenarios more recently.
This includes the modelling of epidemics, especially with the intention of de-
veloping defence strategies. One important class of models (see for example
Nowak and May, 1994) considers the evolution of the virulence of a dis-
ease as the epidemic spreads. An exciting new line of research has recently
been developed which considers the development of cancer as an evolutionary
game, where the population of cancer cells evolves in the environment of the
individual person or animal (Gatenby et al, 2010). A survey of alternative
approaches is considered in Durrett (2014).

1.12 Cross References (to other articles or chapters)

The theoretical foundations for models surveyed in this Chapter are in Vol-
ume I, Chapter Evolutionary games.

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